

Converting from ASCII to Binary

Conversions from ASCII to binary usually start with keyboard entry. If a single key is typed, the conversion occurs when a 30H is subtracted from the number. If more than one key is typed, conversion from ASCII to binary still requires 30H to be subtracted, but there is one additional step. After subtracting 30H, the number is added to the result after the prior result is first multiplied by 10.

The algorithm for converting from ASCII to binary is:

1. Divide
 1. Begin with a binary result of 0.
 2. Subtract 30H from the character typed on the keyboard to convert it to BCD.
 3. Multiply the result by 10, and then add the new BCD digit.
 4. Repeat steps 2 and 3 until the character typed is not an ASCII-coded number.

Example 7-25 illustrates a procedure (READN) used in a program that implements this algorithm. Here, the binary number returns in the AX register as a 16-bit result, which is then stored in memory location TEMP. If a larger result is required, the procedure must be reworked for a 32-bit addition. Each time this procedure is called, it reads a number from the keyboard until any key other than 0 through 9 is typed.

EXAMPLE 7-25

```

;A program that reads one decimal number from the
;keyboard and stores the binary value at TEMP.
;
.MODEL SMALL ;select TINY model
.DATA ;start DATA segment
0000 0000 TEMP DW ? ;define TEMP
0000 .CODE ;start CODE segment
.STARTUP ;start program
0017 E8 0007 CALL READN ;read a number
001A A3 0000 R MOV TEMP,AX ;save it in TEMP
.EXIT ;exit to DOS
;
;The READN procedure reads a decimal number from the
;keyboard and returns its binary value in AX.
;
0021 READN PROC NEAR

0021 53 PUSH BX ;save BX and CX
0022 51 PUSH CX
0023 B9 000A MOV CX,10 ;load 10 for decimal
0026 BB 0000 MOV BX,0 ;clear result
0029 READN1:
0029 B4 01 MOV AH,1 ;read key with echo
002B CD 21 INT 21H

002D 3C 30 CMP AL,'0'
002F 72 14 JB READN2 ;if below '0'
0031 3C 39 CMP AL,'9'
0033 77 10 JA READN2 ;if above '9'

0035 2C 30 SUB AL,'0' ;convert to ASCII

0037 50 PUSH AX ;save digit
0038 8B C3 MOV AX,BX ;multiply result by 10
003A F7 E1 MUL CX
003C 8B D8 MOV BX,AX
003E 58 POP AX
003F B4 00 MOV AH,0
0041 03 D8 ADD BX,AX ;add digit value to result

```

```

0043 EB E4          JMP READN1          ;repeat
0045                READN2:
0045 8B C3          MOV AX,BX          ;get binary result into AX
0047 59            POP CX            ;restore CX and BX
0048 5B            POP BX
0049 C3            RET

004A                READN  ENDP
                        END          ;end of file

```

Displaying and Reading Hexadecimal

Hexadecimal data are easier to read from the keyboard and display than decimal data. These types of data are not used at the applications level, but at the system level. System-level data are often hexadecimal, and must either be displayed in hexadecimal form or read from the keyboard as hexadecimal data.

Reading Hexadecimal Data. Hexadecimal data appear as 0 to 9 and A to F. The ASCII codes obtained from the keyboard for hexadecimal data are 30H to 39H for the numbers 0 through 9, and 41H to 46H (A–F) or 61H to 66H (a–f) for the letters. To be useful, a procedure that reads hexadecimal data must be able to accept both lowercase and uppercase letters.

Example 7–26 shows two procedures: one (CONV) converts the contents of the data in AL from ASCII code to a single hexadecimal digit, and the other (READH) reads a four-digit hexadecimal number from the keyboard and returns with it in register AX. This procedure can be modified to read any-sized hexadecimal number from the keyboard.

EXAMPLE 7–26

```

;A program that reads a 4-digit hexadecimal number from
;the keyboard and stores the result in word-sized
;memory location TEMP.
;
        .MODEL SMALL          ;select SMALL model
0000    .DATA                  ;start DATA segment
0000 0000    TEMP DW ?         ;define TEMP
0000    .CODE                  ;start CODE segment
        .STARTUP              ;start program
0017 E8 0007    CALL READH     ;read hexadecimal number
001A A3 0000 R    MOV TEMP,AX  ;save it at TEMP
        .EXIT                  ;exit to DOS
;
;The READH procedure that reads a 4-digit hexadecimal
;number from the keyboard and returns it in AX.
;This procedure does next check for errors and uses CONV.
;
0021    READH PROC NEAR

0021 51            PUSH CX          ;save BX and CX
0022 53            PUSH BX
0023 B9 0004       MOV CX,4          ;load CX and SI with 4
0026 8B F1         MOV SI,CX
0028 BB 0000       MOV BX,0          ;clear result
002B                READH1:
002B B4 01         MOV AH,1          ;read a key with echo
002D CD 21         INT 21H
002F E8 000A       CALL CONV        ;convert to binary
0032 D3 E3         SHL BX,CL
0034 02 D8         ADD BL,AL        ;form result in BX
0036 4E            DEC SI
0037 75 F2         JNZ READH1       ;repeat 4 times
0039 8B C3         MOV AX,BX
003B 5B            POP BX          ;restore BX and CX

```

```

003C 59                POP CX
003D C3                RET
003E                READH ENDP
;
;The CONV procedure converts AL into hexadecimal.
;
003E                CONV PROC NEAR

003E 3C 39            CMP AL,'9'
0040 76 08            JBE CONV2           ;if 0 through 9
0042 3C 61            CMP AL,'a'
0044 72 02            JB CONV1           ;if uppercase A through F
0046 2C 20            SUB AL,20H        ;convert to uppercase
0048                CONV1:
0048 2C 07            SUB AL,7
004A                CONV2:
004A 2C 30            SUB AL,30H
004C C3                RET
004D                CONV ENDP
;end of file

```

Displaying Hexadecimal Data. To display hexadecimal data, a number must be divided into four-bit segments that are converted into hexadecimal digits. Conversion is accomplished by adding a 30H to the numbers 0 to 9 and a 37H to the letters A to F.

A procedure (DISPH) that displays the contents of the AX register on the video display appears in the program of Example 7-27. Here, the number is rotated left so that the leftmost digit is displayed first. Because AX contains a four-digit hexadecimal number, the procedure displays four hexadecimal digits.

EXAMPLE 7-27

```

;A program that displays the hexadecimal value in AX.
;This program uses DISPH to display a 4-digit value.
;
.MODEL TINY           ;select TINY model
.CODE                ;start CODE segment
.STARTUP             ;start program
0100 B8 0ABC         MOV AX,0ABC         ;load AX with test data
0103 E8 0004         CALL DISPH         ;display AX in hexadecimal
;                   ;exit to DOS
;
;The DISPH procedure displays AX as a 4-digit hex number.
;
010A                DISPH PROC NEAR

010A 53              PUSH BX           ;save BX and CX
010B 51              PUSH CX
010C B1 04           MOV CL,4         ;load rotate count
010E B5 04           MOV CH,4         ;load digit count
0110                DISPH1:
0110 D3 C0           ROL AX,CL         ;position digit
0112 50              PUSH AX
0113 24 0F           AND AL,0FH       ;convert it to ASCII
0115 04 30           ADD AL,30H
0117 3C 39           CMP AL,'9'
0119 76 02           JBE DISPH2
011B 04 07           ADD AL,7
011D                DISPH2:
011D B4 02           MOV AH,2         ;display hexadecimal digit
011F 8A D0           MOV DL,AL
0121 CD 21           INT 21H
0123 58              POP AX
0124 FE CD           DEC CH

```

```

0126 75 E8          JNZ  DISPH1      ;repeat for 4 digits
0128 59            POP  CX          ;restore registers
0129 5B            POP  BX
012A C3            RET

012B              DISPH  ENDP
                          END                ;end of file
    
```

Using Lookup Tables for Data Conversions

Lookup tables are often used to convert data from one form to another. A lookup table is formed in the memory as a list of data that is referenced by a procedure to perform conversions. In the case of many lookup tables, the XLAT instruction can often be used to look up data in a table, provided that the table contains eight-bit wide data and its length is less than or equal to 256 bytes.

Converting from BCD to 7-segment Code. One simple application that uses a lookup table is BCD to 7-segment code conversion. Example 7-28 illustrates a lookup table that contains the 7-segment codes for the numbers 0 to 9. These codes are used with the 7-segment display pictured in Figure 7-1. This 7-segment display uses active high (logic 1) inputs to light a segment. The code is arranged so that the **a** segment is in bit position 0 and the **g** segment is in bit position 6. Bit position 7 is 0 in this example, but it can be used for displaying a decimal point.

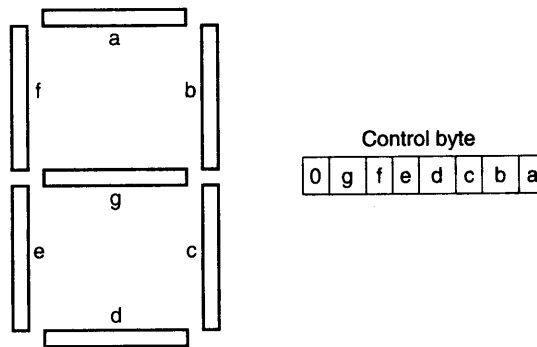


FIGURE 7-1 The 7-segment display.

EXAMPLE 7-28

```

0000          SEG7  PROC  FAR

0000 53          PUSH  BX
0001 BB 0008 R   MOV   BX,OFFSET TABLE
0004 2E: D7     XLAT  CS:TABLE      ;see text
0006 5B          POP   BX
0007 CB          RET

0008 3F          TABLE DB   3FH      ;0
0009 06          DB     6           ;1
000A 5B          DB     5BH        ;2
000B 4F          DB     4FH        ;3
000C 66          DB     66H        ;4
000D 6D          DB     6DH        ;5
000E 7D          DB     7DH        ;6
000F 07          DB     7           ;7
0010 7F          DB     7FH        ;8
0011 6F          DB     6FH        ;9

0012          SEG7  ENDP
    
```

The procedure that performs the conversion contains only two instructions and assumes that AL contains the BCD digit to be converted to 7-segment code. One of the instructions addresses the lookup table by loading its address into BX, and the other performs the conversion and returns the 7-segment code in AL.

Because the lookup table is located in the code segment and the XLAT instruction accesses the data segment by default, the XLAT instruction includes a segment override. Notice that a dummy operand (TABLE) is added to the XLAT instruction so that the (CS:) code segment override prefix can be added to the instruction. Normally,

XLAT does not contain an operand unless its default segment must be overridden. The LODS and MOVS instructions are also overridden in the same manner as XLAT by using a dummy operand.

Using a Lookup Table to Access ASCII Data. Some programming techniques require that numeric codes be converted to ASCII character strings. For example, suppose that you need to display the days of the week for a calendar program. Because the number of ASCII characters in each day is different, some type of lookup table must be used to reference the ASCII-coded days of the week.

The program in Example 7-29 shows a table that references ASCII-coded character strings located in the code segment. Each character string contains an ASCII-coded day of the week. The table references each day of the week. The procedure that accesses the day of the week uses the AL register, and the numbers 0 to 6 to refer to Sunday through Saturday. If AL contains a 2 when this procedure is called, the word "Tuesday" is displayed on the video screen.

EXAMPLE 7-29

```

;A program that displays the current day of the
;week by using the system clock/calendar.
;
.MODEL SMALL ;select SMALL model
.DATA ;start DATA segment
0000 DTAB DW SUN, MON, TUE, WED, THU, FRI, SAT
0000 000E R 0015 R
001C R 0024 R
002E R 0037 R
003E R
000E 53 75 6E 64 61 79 SUN DB 'Sunday$'
24
0015 4D 6F 6E 64 61 79 MON DB 'Monday$'
24
001C 54 75 65 73 64 61 TUE DB 'Tuesday$'
79 24
0024 57 65 64 6E 65 73 WED DB 'Wednesday$'
64 61 79 24
002E 54 68 75 72 73 64 THU DB 'Thursday$'
61 79 24
0037 46 72 69 64 61 79 FRI DB 'Friday$'
24
003E 53 61 74 75 72 64 SAT DB 'Saturday$'
61 79 24
0000 .CODE ;start CODE segment
.STARTUP ;start program
0017 B4 2A MOV AH, 2AH ;get day of week
0019 CD 21 INT 21H ;access DOS
001B E8 0004 CALL DAYS ;display day of week
.EXIT ;exit to DOS
0022 DAYS PROC NEAR
0022 52 PUSH DX ;save DX and SI
0023 56 PUSH SI
0024 BE 0000 R MOV SI, OFFSET DTAB ;address table
0027 B4 00 MOV AH, 0 ;find day of week
0029 03 C0 ADD AX, AX
002B 03 F0 ADD SI, AX
002D 8B 14 MOV DX, [SI] ;get day of week
002F B4 09 MOV AH, 9 ;display string
0031 CD 21 INT 21H
0033 5E POP SI ;restore registers
0034 5A POP DX
0035 C3 RET

```

```
0036          DAYS      ENDP
                   END          ;end of file
```

This procedure first addresses the table by loading its address into the SI register. Next, the number in AL is converted into a 16-bit number and doubled because the table contains two bytes for each entry. This index is then added to SI to address the correct entry in the lookup table. The address of the ASCII character string is now loaded into DX by the MOV DX,CS:[SI] instruction.

Before the INT 21H DOS function is called, the DS register is placed on the stack and loaded with the segment address of CS. This allows DOS function number 09H (display a string) to be used to display the day of the week. This procedure converts the numbers 0 to 6 to the days of the week.

An Example Program Using Data Conversions

A program example is required to combine some of the data-conversion DOS functions. Suppose that you must display the time and date on the video screen. This example program (see Example 7-30) displays the time as 10:45 A.M. and the date as Tuesday, May 14, 2002. The program is short because it calls a procedure that displays the time and a second procedure that displays the date.

The time is available from DOS, using an INT 21H function call number 2CH. This returns with the hours in CH and minutes in CL. Also available are seconds in DH and hundredths of seconds in DL. The date is available by using INT 21H function call number 2AH. This leaves the day of the week in AL, the year in CX, the day of the month in DH, and the month in DL.

EXAMPLE 7-30

```

;A program that displays the time and date in the
;form: 10:45 A.M., Tuesday May 14, 2002.
;
        .MODEL SMALL          ;select SMALL model
        .NOLISTMACRO          ;don't expand macros
        .DATA                  ;start CODE segment
0000    DTAB    DW    SUN,MON,TUE,WED,THU,FRI,SAT
0000    0026 R 002F R
          0038 R 0042 R
          004E R 0059 R
          0062 R
000E    MTAB    DW    JAN,FEB,MAR,APR,MAY,JUN
          006D R 0076 R
          0080 R 0087 R
          008E R 0093 R
001A    DW    JUL,AUG,SEP,OCT,NOV,DCE
          0099 R 009F R
          00A7 R 00B2 R
          00BB R 00C5 R
0026    53 75 6E 64 61 79 SUN    DB    'Sunday, $'
          2C 20 24
002F    4D 6F 6E 64 61 79 MON    DB    'Monday, $'
          2C 20 24
0038    54 75 65 73 64 61 TUE    DB    'Tuesday, $'
          79 2C 20 24
0042    57 65 64 6E 65 73 WED    DB    'Wednesday, $'
          64 61 79 2C 20 24
004E    54 68 75 72 73 64 THU    DB    'Thursday, $'
          61 79 2C 20 24
0059    46 72 69 64 61 79 FRI    DB    'Friday, $'
          2C 20 24
0062    53 61 74 75 72 64 SAT    DB    'Saturday, $'
          61 79 2C 20 24
006D    4A 61 6E 75 61 72 JAN    DB    'January $'
          79 20 24
0076    46 65 62 72 75 61 FEB    DB    'February $'
          72 79 20 24
0080    4D 61 72 63 68 20 MAR    DB    'March $'
```

```

24
0087 41 70 72 69 6C 20 APR DB 'April $'
24
008E 4D 61 79 20 24 MAY DB 'May $'
0093 4A 75 6E 65 20 24 JUN DB 'June $'
0099 4A 75 6C 79 20 24 JUL DB 'July $'
009F 41 75 67 75 73 74 AUG DB 'August $'
20 24
00A7 53 65 70 74 65 6D SEP DB 'September $'
62 65 72 20 24
00B2 4F 63 74 6F 62 65 OCT DB 'October $'
72 20 24
00BB 4E 6F 76 65 6D 62 NOV DB 'November $'
65 72 20 24
00C5 44 65 63 65 6D 62 DCE DB 'December $'
65 72 20 24
0000
DISP .CODE ;start CODE segment
MACRO CHAR
PUSH AX ;;save AX and DX
PUSH DX
MOV DL,CHAR ;;display character
MOV AH,2
INT 21H
POP DX ;;restore AX and DX
POP AX
ENDM
0017 E8 0007 .STARTUP ;start program
001A E8 00A3 CALL TIMES ;display time
CALL DATES ;display date
.EXIT ;exit to DOS

0021 TIMES PROC NEAR
0021 B4 2C MOV AH,2CH ;get time from DOS
0023 CD 21 INT 21H
0025 B7 41 MOV BH,'A' ;set 'A' for AM
0027 80 FD 0C CMP CH,12
002A 72 05 JB TIMES1 ;if below 12:00 noon
002C B7 50 MOV BH,'P' ;set 'P' for PM
002E 80 ED 0C SUB CH,12 ;adjust to 12 hours
0031 TIMES1:
0031 0A ED OR CH,CH ;test for 0 hour
0033 75 02 JNE TIMES2 ;if not 0 hour
0035 B5 0C MOV CH,12 ;change 0 hour to 12
0037 TIMES2:
0037 8A C5 MOV AL,CH
0039 B4 00 MOV AH,0
003B D4 0A AAM ;convert hours
003D 0A E4 OR AH,AH
003F 74 0D JZ TIMES3 ;if no tens of hours
0041 80 C4 30 ADD AH,'0' ;convert tens
DISP AH ;display tens
004E TIMES3:
004E 04 30 ADD AL,'0' ;convert units
DISP AL ;display units
DISP ':' ;display colon
MOV AL,CL
MOV AH,0
0064 8A C1 AAM ;convert minutes
0066 B4 00 ADD AX,3030H
0068 D4 0A PUSH AX
006A 05 3030 DISP AH ;display tens
006D 50

```

208 CHAPTER 7 PROGRAMMING THE MICROPROCESSOR

```

0078 58          POP AX
                DISP AL          ;display units
                DISP ' '        ;display space
                DISP BH          ;display 'A' or 'P'
                DISP '.'         ;display .
                DISP 'M'        ;display M
                DISP '.'         ;display .
                DISP ' '        ;display space
00BF C3          RET

00C0          TIMES ENDP

00C0          DATES PROC NEAR

00C0 B4 2A      MOV AH,2AH      ;get date from DOS
00C2 CD 21      INT 21H
00C4 52         PUSH DX
00C5 B4 00      MOV AH,0        ;get day of week
00C7 03 C0      ADD AX,AX
00C9 BE 0000 R  MOV SI,OFFSET DTAB ;address day table
00CC 03 F0      ADD SI,AX
00CE 8B 14      MOV DX,[SI]      ;address day of week
00D0 B4 09      MOV AH,9        ;display day of week
00D2 CD 21      INT 21H
00D4 5A         POP DX
00D5 52         PUSH DX
00D6 8A C6      MOV AL,DH        ;get month
00D8 FE C8      DEC AL
00DA B4 00      MOV AH,0
00DC 03 C0      ADD AX,AX
00DE BE 000E R  MOV SI,OFFSET MTAB ;address month table
00E1 03 F0      ADD SI,AX
00E3 8B 14      MOV DX,[SI]      ;address month
00E5 B4 09      MOV AH,9        ;display month
00E7 CD 21      INT 21H
00E9 5A         POP DX
00EA 8A C2      MOV AL,DL        ;get day of month
00EC B4 00      MOV AH,0
00EE D4 0A      AAM            ;convert to BCD
00F0 0A E4      OR AH,AH
00F2 74 0D      JZ DATES1       ;if tens is 0
00F4 80 C4 30   ADD AH,30H      ;convert tens
                DISP AH      ;display tens

0101          DATES1:
0101 04 30      ADD AL,30H      ;convert units
                DISP AL      ;display units
                DISP ','      ;display comma
                DISP ' '      ;display space

0121 81 F9 07D0 CMP CX,2000     ;test for year 2000
0125 72 19      JB DATES2       ;if below year 2000
0127 83 E9 64   SUB CX,100      ;scale to 1900 - 1999
                DISP '2'      ;display 2
                DISP '0'      ;display 0

013E EB 14      JMP DATES3

0140          DATES2:
                DISP '1'      ;display 1
                DISP '9'      ;display 9

0154          DATES3:
0154 81 E9 076C SUB CX,1900     ;scale to 00 - 99
0158 8B C1      MOV AX,CX
015A D4 0A      AAM            ;convert to BCD

```



```

015C 05 3030          ADD  AX,3030H          ;convert to ASCII
                                DISP AH          ;display tens
                                DISP AL          ;display units
0173 C3              RET
0174                DATES  ENDP
                                END              ;end of file

```

This procedure uses two ASCII lookup tables that convert the day and month to ASCII character strings. It also uses the AAM instruction to convert from binary to BCD for the time and date. The displaying of data is handled in two ways: by character string (function 09H) and by single character (function 06H).

The memory model (SMALL) consists of two segments: .DATA and .CODE. The data segment contains the character strings used with the procedures that display time and date. The code segment contains TIMES and DATES procedures, and a macro (DISP) that displays an ASCII character. The main program is very short and consists of two CALL instructions. The year 2000 problem is corrected in this program, but not the year 2100 problem.

Numeric Sort Program

At times, numbers must be sorted into numeric order. This is often accomplished with a bubble sort. Figure 7-2 shows five numbers that are sorted with a bubble sort. Notice that the set of five numbers is tested four times with four passes. For each pass, two consecutive numbers are compared and sometimes exchanged. Also notice that during the first pass, there are four comparisons, during the second three, etc.

Example 7-31 illustrates a program that accepts 10 numbers from the keyboard (0-65535). After these 16-bit numbers are accepted and stored in memory section ARRAY, they are sorted by using the bubble-sorting technique. This bubble sort uses a flag to determine whether any numbers were exchanged in a pass. If no numbers were exchanged, the numbers are in order and the sort terminates.

EXAMPLE 7-31

```

                                .MODEL SMALL
                                .DATA
0000 000A {                      ARRAY DW 10 DUP (?)
;array
                                0000
                                }
0014 0D 0A 45 6E 74 65          MES1 DB 13,10,'Enter 10 numbers:',13,10,10,'$'
                                72 20 31 30 20 6E
                                75 6D 62 65 72 73
                                3A 0D 0A 0A 24
002B 0D 0A 0A 53 6F 72          MES2 DB 13,10,10,'Sorted Data:',13,10,10,'$'
                                74 65 64 20 44 61
                                74 61 3A 0D 0A 0A
                                24
0000                                .CODE
                                DISP  MACRO  PARA
                                PUSH  AX
                                MOV   AH,6
                                MOV   DL, PARA
                                INT   21H

```

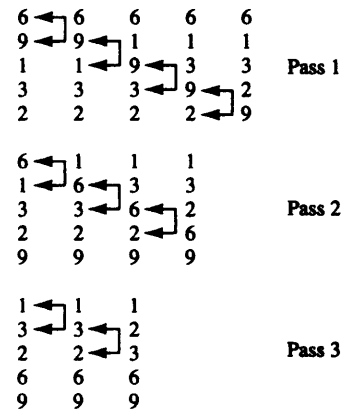


FIGURE 7-2 A bubble sort showing data as they are sorted. Note: Sorting five numbers may require four passes.

210 CHAPTER 7 PROGRAMMING THE MICROPROCESSOR

```

        POP    AX

        ENDM
        GET    MACRO

        .REPEAT
            MOV    AH,6
            MOV    DL,-1
            INT    21H
        .UNTIL (AL>='0' && AL <='9') || AL==13 || AL==','
        DISP    AL
        .IF AL==13
            DISP    10
        .ENDIF
        .IF AL>='0' && AL<='9'
            SUB    AL,'0'
        .ENDIF

        ENDM

        STRING    MACRO WHERE

            MOV    DX,OFFSET WHERE
            MOV    AH,9
            INT    21H

        ENDM

        .STARTUP

        STRING    MES1
        CLD
        MOV    CX,10
        MOV    DI,OFFSET ARRAY
        MOV    AX,DS
        MOV    ES,AX
        .REPEAT
            CALL    GETN    ;get 10 numbers
        .UNTILCXZ
        STRING    MES2
        CALL    SORT    ;sort 10 numbers
        MOV    CX,9
        MOV    SI,OFFSET ARRAY
        .REPEAT
            CALL    DISPN    ;display 10 numbers
            DISP    ','
        .UNTILCXZ
        CALL    DISPN
        .EXIT

        0052                GETN    PROC    NEAR

        0052    BD 000A        MOV    BP,10
        0055    BB 0000        MOV    BX,0
                                .WHILE 1
                                GET
                                .BREAK .IF AL==13 || AL==','
        0094    93                XCHG    .AX,BX
        0095    F7 E5            MUL    BP
        0097    93                XCHG    AX,BX
        0098    B4 00            MOV    AH,0
        009A    03 D8            ADD    BX,AX
                                .ENDW
        009E    8B C3            MOV    AX,BX
        00A0    AB                STOSW

```

```

00A1  C3                                RET
00A2                                GETN  ENDP
00A2                                DISPN PROC    NEAR
00A2  BB 000A                          MOV    BX, 10
00A5  53                                PUSH   BX
00A6  AD                                LODSW
                                .REPEAT
00A7  BA 0000                          MOV    DX, 0
00AA  F7 F3                            DIV    BX
00AC  52                                PUSH   DX
                                .UNTIL AX==0
                                .WHILE 1
00B1  58                                POP    AX
                                .BREAK .IF AL==10
00B6  04 30                            ADD    AL, '0'
                                DISP   AL
                                .ENDW
00C2  C3                                RET
00C3                                DISPN  ENDP
00C3                                SORT   PROC    NEAR
00C3  BB 0009                          MOV    BX, 9
                                .REPEAT
00C6  8B CB                            MOV    CX, BX
00C8  BE 0000 R                        MOV    SI, OFFSET ARRAY
00CB  B2 00                            MOV    DL, 0
                                .REPEAT
00CD  AD                                LODSW
00CE  3B 04                            CMP    AX, [SI]
                                .IF !CARRY?
00D2  8B 2C                            MOV    BP, [SI]
00D4  89 6C FE                        MOV    [SI-2], BP
00D7  89 04                            MOV    [SI], AX
00D9  FE C2                            INC    DL
                                .ENDIF
                                .UNTILCXZ
00DD  4B                                DEC    BX
                                .UNTIL BX==0 || DL==0
00E6  C3                                RET
00E7                                SORT   ENDP
                                END

```

Once the numbers are sorted, they are displayed on the video screen in ascending numerical order. No provision is made for errors as each number is typed. The program terminates after sorting one set of 10 numbers and must be invoked again to sort 10 new numbers.

7-4 INTERRUPT HOOKS

Hooks are used to tap into or intercept the interrupt structure of the microprocessor. For example, we might hook into the keyboard interrupt so that we can detect a special keystroke called a *hot key*. Whenever the hot key is typed, we can access a terminate and stay resident (TSR) program that performs a special task. Some examples of hot key applications are pop-up calculators and pop-up clocks.

Intercepting an Interrupt

In order to intercept an interrupt, we must use a DOS function call that reads the current address from the interrupt vector. DOS function call number 35H is used to read the current interrupt vector and DOS function call number 25H is used to change the address of the current vector. In both DOS function calls, AL indicates the vector type number (00H–FFH) and AH indicates the DOS function call number.

When the vector is read by using function 35H, the offset address is returned in register BX and the segment address is in register ES. These two registers are saved so that they can be restored when the interrupt hook is removed from memory. When the vector is set, it is set to the address stored at the memory location addressed by DS:DX.

The process of installing an interrupt handler through a hook is illustrated in the program of Example 7–32. This program intercepts the divide error interrupt by first reading the current interrupt vector address and storing it into a double-word memory location for access by the new interrupt service procedure. Next, the address of the new interrupt service procedure, stored in DS:DX, is placed into the vector using DOS function call number 25H.

EXAMPLE 7–32

```

;A sequence of instructions that show the installation
;or a new interrupt for vector 0 (divide error).
;Note this is not a complete program.
;
                                .MODEL TINY
0000                                .CODE
                                .STARTUP
0100 EB 05                        JMP MAIN                        ;skip
0102 00000000 ADDR DD ?          ;old interrupt vector
0106                                NEW PROC FAR                       ;new interrupt procedure
0106 CF                            IRET                                ;do nothing interrupt
0107                                NEW ENDP
0107                                MAIN:
0107 8C C8                            MOV AX,CS                      ;address CS with DS
0109 8E D8                            MOV DS,AX
                                ;get vector 0 address
010B B8 3500                            MOV AX,3500H
010E CD 21                            INT 21H
                                ;save vector address at ADDR
0110 89 1E 0102 R                        MOV WORD PTR ADDRESS,BX
0114 8C 06 0104 R                        MOV WORD PTR ADDRESS+2,ES
                                ;install new interrupt vector 0 address
0118 B8 2500                            MOV AX,2500H
011B BA 0106 R                            MOV DX,OFFSET NEW
011E CD 21                            INT 21H
                                ;other installation software continues here

```

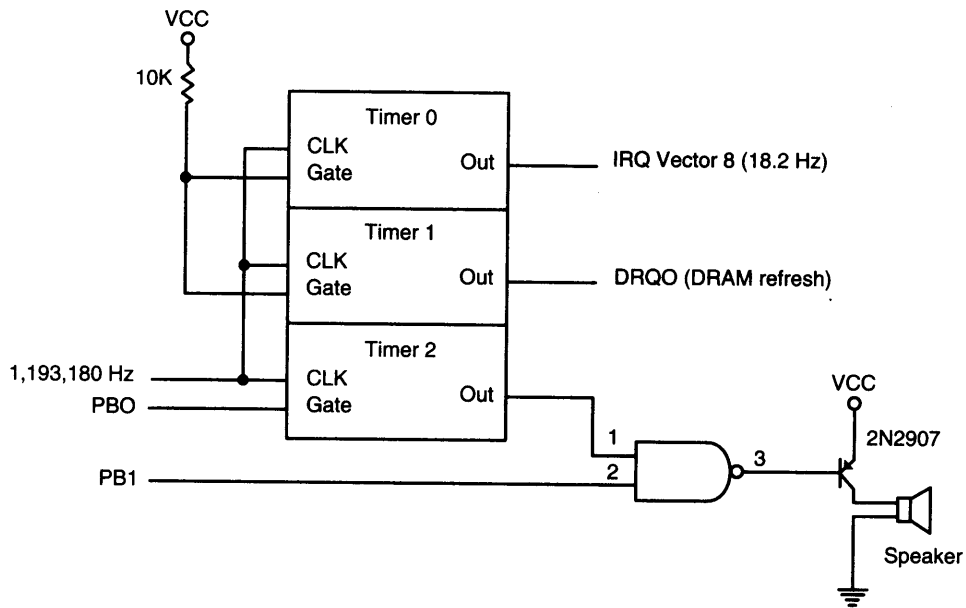


FIGURE 7-3 The speaker and timer circuit in the personal computer (I/O ports 40–43H program the timer, and I/O port 61H programs PB0 and PB1).

Example TSR Alarm

A simple example showing an interrupt hook and TSR causes a beep on the speaker after one hour or one-half hour. We all seem to get lost in computer processing, and this program makes it easy to keep track of time because of the audible beep.

The beep is caused by using timer 2 of the timer found inside the PC in order to generate an audio tone at the speaker. (See Section 12-5 for a discussion of the timer and see Figure 7-3 for its connection in the computer.) Programming timer 2 with a particular beep frequency or tone is accomplished by programming timer 2 with 1,193,180, divided by the desired tone. For example, if we divide 1,193,180 by 800, the speaker generates an 800 Hz audio tone. See the BEEP procedure (shown in Example 7-33) for programming the timer, and turning the speaker on and off after a short wait determined by the number of clock ticks. This procedure uses six clock ticks to produce a beep lasting $\frac{1}{3}$ second. Note that each clock tick occurs about 18.2 times a second (the actual time is closer to 18.206). This is accomplished by using the user wait timer locations in the first segment of the memory. The user wait timer is updated 18.2 times per second by the computer so that it can be used to time events. The program that uses the BEEP procedure causes an audio tone of 1000 Hz, 1200 Hz, and 1400 Hz (each with a $\frac{1}{3}$ -second duration) to repeat four times.

EXAMPLE 7-33

```

;A program that beeps the speaker with some sample audio
;tones that each have a duration of 1/3 second.
;
.MODEL TINY
.CODE
.STARTUP
0100 B8 0000    MOV AX,0
0103 8E D8     MOV DS,AX           ;address segment 0000H
0105 B9 0004    MOV CX,4           ;set count to 4
0108 E4 61     IN AL,61H         ;enable timer and speaker

```

214 CHAPTER 7 PROGRAMMING THE MICROPROCESSOR

```

010A 0C 03          OR   AL,3           ;set PB0 and PB1
010C E6 61          OUT  61H,AL
010E                MAIN1:
010E BB 03E8        MOV  BX,1000        ;select 1000 Hz tone
0111 E8 0018        CALL BEEP
0114 BB 04B0        MOV  BX,1200        ;select 1200 Hz tone
0117 E8 0012        CALL BEEP
011A BB 0578        MOV  BX,1400        ;select 1400 Hz tone
011D E8 000C        CALL BEEP
0120 E2 EC          LOOP MAIN1        ;repeat 4 times

0122 E4 61          IN   AL,61H          ;turn speaker off
0124 34 03          XOR  AL,3           ;clear PB0 and PB1
0126 E6 61          OUT  61H,AL
                .EXIT

;
;The BEEP procedure programs timer 2 to beep the speaker
;for 1/3 of a second with the frequency BX.
;***input parameters***
;BX = desired audio tone
;***uses***
;WAITS procedure to wait for 1/3 second
;
012C                BEEP  PROC NEAR        ;beep speaker 1/3 second

012C B8 34DC        MOV  AX,34DCH        ;load AX with 1,193,180
012F BA 0012        MOV  DX,12H
0132 F7 F3          DIV  BX           ;find count
0134 E6 42          OUT  42H,AL        ;program timer 2
0136 8A C4          MOV  AL,AH
0138 E6 42          OUT  42H,AL
013A E8 0001        CALL WAITS        ;wait 1/3 second
013D C3            RET

013E                BEEP  ENDP

;
;the WAITS procedure waits 1/3 of a second
;***uses***
;memory doubleword location 0000:46CH to time the wait
;
013E                WAITS PROC NEAR

013E BA 0006        MOV  DX,6           ;number of clock ticks
0141 BB 0000        MOV  BX,0
0144 03 16 046C    ADD  DX,DS:[46CH]   ;get tick count plus time
0148 13 1E 046E    ADC  BX,DS:[46EH]
014C                WAIT1:
014C 8B 2E 046C    MOV  BP,DS:[46CH]   ;test for elapsed time
0150 A1 046E        MOV  AX,DS:[46EH]
0153 2B EA          SUB  BP,DX
0155 1B C3          SBB  AX,BX
0157 72 F3          JC   WAIT1        ;keep testing

0159 C3            RET

015A                WAITS ENDP
                END

```

The CHIME program (see Example 7-34) hooks into interrupt vector 8 and beeps the speaker once each half-hour and twice on the hour. This program is a TSR and remains active until the computer is turned off. Note how the TSR is installed and how the interrupt vector is hooked. Also notice that the normal interrupt vector 8 procedure continues to execute, even as the beeper is activated.

EXAMPLE 7-34

```

;A terminate and stay resident program that hooks into
;interrupt vector 8 to beep the speaker one time per
;half-hour and two times per hour.
;***must be assembled as a .COM file*** for use with
;version 5.10 of MASM
        .MODEL TINY
0000        .CODE
        .STARTUP
0100  E9 00CE        JMP  INSTALL        ;install interrupt
        = 03E8        TONE  EQU  1000        ;set tone at 1000 Hz
0103  00            COUNT DB  0            ;elapsed time counter
0104  00000000      ADD8  DD  ?            ;old vector address
0108  00            PASS  DB  0            ;1 or 2 beeps
0109  00            BEEP  DB  0            ;beep or silent
010A  00            FLAG  DB  0            ;busy flag

010B                VEC8  PROC FAR        ;interrupt procedure

010B  2E: 80 3E 010A R    CMP  CS:FLAG,0        ;test busy flag
        00
0111  74 05            JE   VEC81        ;if not busy
0113  2E: FF 2E 0104 R    JMP  CS:ADD8        ;if busy do normal INT 8
0118                VEC81:
0118  9C            PUSHF        ;do normal INT 8
0119  2E: FF 1E 0104 R    CALL CS:ADD8
011E  2E: C6 06 010A R    MOV  CS:FLAG,1        ;show busy
        01
0124  FB            STI            ;allow other interrupts
0125  2E: 80 3E 0108 R    CMP  CS:PASS,0
        00
012B  75 2C            JNE  VEC83        ;if beep counter active
012D  50            PUSH  AX        ;save registers
012E  51            PUSH  CX
012F  52            PUSH  DX
0130  B4 02            MOV  AH,2        ;get time from BIOS
0132  CD 1A            INT  1AH
0134  80 FE 00        CMP  DH,0        ;is it 00 seconds
0137  75 68            JNE  VEC86        ;not time yet, so return
0139  80 F9 00        CMP  CL,0        ;test for hour
013C  74 10            JE   VEC82        ;if hour beep 2 times
013E  80 F9 30        CMP  CL,30H      ;test for half-hour
0141  75 5E            JNE  VEC86        ;if not half-hour
0143  E8 0065        CALL BEEPS        ;start speaker beep
0146  2E: C6 06 0108 R    MOV  CS:PASS,1        ;set number of beeps to 1
        01
014C  EB 53            JMP  VEC86        ;end it
014E                VEC82:
014E  E8 005A        CALL BEEPS        ;start speaker beep
0151  2E: C6 06 0108 R    MOV  CS:PASS,2        ;set number of beeps to 2
        02
0157  EB 48            JMP  VEC86        ;end it
0159                VEC83:
0159  2E: 80 3E 0103 R    CMP  CS:COUNT,0    ;test for end of delay
        00
015F  74 07            JE   VEC84        ;if time delay has elapsed
0161  2E: FE 0E 0103 R    DEC  CS:COUNT
0166  EB 3C            JMP  VEC88        ;end it
0168                VEC84:
0168  2E: 80 3E 0109 R    CMP  CS:BEEP,0      ;test beep on
        00
016E  75 1C            JNE  VEC85        ;if beep is on
0170  2E: FE 0E 0108 R    DEC  CS:PASS        ;test for 2 beeps

```

216 CHAPTER 7 PROGRAMMING THE MICROPROCESSOR

```

0175 74 2D          JZ   VEC88          ;if second beep not needed
0177 2E: C6 06 0103 R  MOV  CS:COUNT,9      ;reset count
      09
017D 2E: C6 06 0109 R  MOV  CS:BEEP,1          ;beep on for second beep
      01
0183 50            PUSH AX
0184 E4 61          IN   AL,61H          ;enable speaker for beep
0186 0C 03          OR   AL,3
0188 E6 61          OUT  61H,AL
018A EB 17          JMP  VEC87          ;end it
018C                VEC85:
018C 2E: C6 06 0103 R  MOV  CS:COUNT,9      ;reset count
      09
0192 2E: C6 06 0109 R  MOV  CS:BEEP,0          ;show beep is off
      00
0198 50            PUSH AX
0199 E4 61          IN   AL,61H          ;disable speaker
019B 34 03          XOR  AL,3
019D E6 61          OUT  61H,AL
019F EB 02          JMP  VEC87          ;end it
01A1                VEC86:
01A1 5A            POP  DX          ;restore registers
01A2 59            POP  CX
01A3                VEC87:
01A3 58            POP  AX
01A4                VEC88:
01A4 2E: C6 06 010A `  MOV  CS:FLAG,0        ;show not busy
      00
01AA CF            IRET          ;interrupt return

01AB                VEC8  ENDP
;
;The BEEPS procedure programs the speaker for the
;frequency stored as TONE using an equate at assembly
;time. The duration of the beep is 1/2 second.
;***uses registers AX, CX, and DX***

01AB                BEEPS  PROC NEAR          ;beep speaker

01AB 2E: 8B 0E 03E8      MOV  CX,CS:TONE      ;set tone
01B0 B8 34DC          MOV  AX,34DCH        ;load AX with 1,193,180
01B3 BA 0012          MOV  DX,12H
01B6 F7 F1            DIV  CX              ;calculate count
01B8 E6 42            OUT  42H,AL          ;program timer 2
01BA 8A C4            MOV  AL,AH
01BC E6 42            OUT  42H,AL

01BE E4 61            IN   AL,61H          ;speaker on
01C0 0C 03            OR   AL,3
01C2 E6 61            OUT  61H,AL
01C4 2E: C6 06 0103 R  MOV  CS:COUNT,9      ;set count for 1/2 second
      09
01CA 2E: C6 06 0109 R  MOV  CS:BEEP,1          ;indicate beep is on
      01
01D0 C3              RET

01D1                BEEPS  ENDP

01D1                INSTALL:
01D1 8C C8            MOV  AX,CS          ;install interrupt VEC8
01D3 8E D8            MOV  DS,AX          ;overlap CS and DS

01D5 B8 3508          MOV  AX,3508H        ;get current vector 8
01D8 CD 21            INT  21H            ;and save it

```



```

01DA 89 1E 0104 R      MOV  WORD PTR ADD8,BX
01DE 8C 06 0106 R      MOV  WORD PTR ADD8+2,ES

01E2 B8 2508          MOV  AX,2508H
01E5 BA 010B R      MOV  DX,OFFSET VEC8      ;address interrupt VEC8
01E8 CD 21          INT  21H                ;install vector 8

01EA BA 01D1 R      MOV  DX,OFFSET INSTALL  ;find paragraphs
01ED B1 04          MOV  CL,4
01EF D3 EA          SHR  DX,CL
01F1 42          INC  DX

01F2 B8 3100          MOV  AX,3100H           ;exit to DOS as TSR
01F5 CD 21          INT  21H
END

```

The CHIME program uses several memory locations as flags to signal the operation of the interrupt service procedure. The first flag tested by CHIME is the busy flag (FLAG), which indicates that a part of the interrupt service procedure is active. If FLAG = 1 (busy condition), the procedure jumps to the normal vector 8 interrupt (JMP CS:ADD8), which ends VEC8's execution. If FLAG = 0 (not busy), the interrupt service procedure continues at VEC81. The default address for all direct memory data is the data segment. In the TSR software used in this example and others, it is important to use the segment override prefix (CS:) to ensure that the program addresses data in the code segment, where it appears.

At VEC81, the normal vector 8 interrupt is executed with a forced interrupt call (PUSHF followed by a CALL CS:ADD8). Upon return from the normal vector 8 interrupt (required to keep accurate time), the busy flag is set to show a busy condition (FLAG = 1) and other interrupts are enabled with the STI instruction.

The PASS flag is now tested to see if the VEC8 procedure is currently beeping the speaker. If PASS = 0 (not beeping speaker), the time of day is retrieved from BIOS by using the INT 1AH instruction. It is important not to access DOS from within a TSR or interrupt service procedure. If DOS is accessed at this time, it may be in the process of executing an operation that affects the interrupt. This would cause the program to crash. The INT 1AH instruction returns the number of seconds (DH), minutes (CL), and hours (CH) in BCD form. After obtaining the current time, the number of seconds is tested for zero. If it is not zero seconds, the interrupt procedure ends. If it is zero seconds, then CL is tested for 00 minute (hour) and 30 minutes. If either case is true, the speaker is enabled and TONE is programmed in the timer by a call to BEEPS. If neither case is true, the interrupt ends. Notice that the BEEPS procedure programs timer 2, enables the speaker, and sets the count to 9.

The time delay counter (COUNT) is decremented each time the interrupt occurs. If the count reaches zero, the procedure tests BEEP to control the speaker. If the speaker is beeping, the procedure turns it off and resets the time delay count to 9. If the speaker is not beeping, the procedure tests PASS to determine if another beep is required on the hour. The time delay is $\frac{1}{2}$ second (COUNT = 9) in this program and cannot be less. If a delay of less than $\frac{1}{2}$ second is chosen, the speaker will beep twice, for both the hour and half hour. The reason is that the clock (INT 1AH) is checked for the zero second. If a time delay of less than $\frac{1}{2}$ second is used, the half-hour will be picked up twice.

The TSR program is loaded into memory at the DOS command line by typing the name of the program; in this case, the program is called CHIME. If DOS version 5.0 or 6.X is in use, you can load CHIME into the upper memory or high memory area by typing LOADHIGH CHIME. Once this program loads into memory, it remains in the background, beeping off the time until the power to the computer is disconnected or until it is rebooted. This is an excellent, not too annoying addition to the system to keep track of time. The next section of the text describes hot keys. If desired, a hot key could be used to enable and disable CHIME.

Example Hot Key Program

Hot keys are keystrokes that invoke terminate and stay resident programs. For example, an ALT + C key could be defined as a hot key that calls a program that displays the time. Note that the hot key is detected inside most applications, but not at the DOS command line, where it may lock up the system if used. To detect a hot key, we usu-

ally hook into interrupt vector 9, which is the keyboard interrupt that occurs if any key is typed. This allows us to test the keyboard and detect a hot key before the normal interrupt processes the keystroke.

A hot key is installed with a TSR program and an interrupt hook. To illustrate a hot key program that can be useful, a program is developed that counts keystrokes. The keystroke counter program (see Example 7-35) is useful in a business environment that uses computers for data entry or other tasks. With this type of program, productivity can be assessed. The keystroke counter program counts each keystroke and only displays the count when the ALT + K key is pressed. (It is important to note that this program spies on workers and it is the duty of any company using the program to notify the worker. It may even be the responsibility of the company to obtain permission from the worker before a program such as this is placed into service.)

This program can be modified to keep track of keystrokes by the hour or any other time unit. In this example, the keystroke count (up to four billion) accumulates keystrokes for as long as power is applied to the computer. The program also stores the installation time for security purposes. This is important because if a machine is reset, the start time for this TSR will be reset.

This program hooks into interrupt 8 and 9 to count keys. The interrupt 9 hook detects the hot key (ALT + K) and counts keystrokes. When the hot key is detected, the 18.2 Hz interrupt 8 activates the hot key program that displays the keystroke count and time of installation. This type of TSR is often called a *pop-up program* because it pops up when the hot key is typed. Notice that this program uses INT 16H to test the keyboard. Never use a DOS INT 21H function call within a TSR or interrupt hook because serious problems can arise. This program also uses direct manipulation of the video text memory that begins at location B8000H. This memory is organized with two bytes per ASCII character. The first byte contains the ASCII code, and the following byte contains the background and character color.

EXAMPLE 7-35

```

;A TSR program that counts keystrokes and reports the
;time of installation and number of accumulated
;keystrokes when the ALT-K key combination is activated.
;***requires an 80386 or newer microprocessor***
;XX for use on PC XXX
.MODEL TINY
.386
.0000 .CODE
.STARTUP
0100 E9 0241 JMP INSTALL ;install VEC8 and VEC9

0103 00 HFLAG DB 0 ;Hot-key detected
0104 00000000 ADD8 DD ? ;old vector 8 address
0108 00000000 ADD9 DD ? ;old vector 9 address
010C 00000000 COUNT DD 0 ;Keystroke counter
0110 00 HOUR DB ? ;start-up time
0111 00 MIN DB ?
0112 0 0 SFLAG DB 0 ;start-up flag
0113 00 FLAG8 DB 0 ;interrupt 8 busy
0114 25 KEY DB 25H ;scan code for K
0115 08 HMASK DB 8 ;alternate key mask
0116 08 MKEY DB 8 ;alternate key
0117 00A0 [ SCRN DB 160 DUP (?) ;screen buffer
00
]
01B7 54 69 6D 65 MES1 DB 'Time = '
20 3D 20
01BE 20 20 20 4B MES2 DB ' KeyStrokes = '
65 79 53 74
72 6F 6B 65
73 20 3D 20

01CE VEC9 PROC FAR ;keyboard intercept

```

```

01CE FB STI ;enable interrupts
01CF 66| 50 PUSH EAX ;save EAX
01D1 E4 60 IN AL,60H ;get scan code
01D3 2E: 3A 06 0114 R CMP AL,CS:KEY ;test for K
01D8 75 16 JNE VEC91 ;no hot-key
01DA B8 0000 MOV AX,0 ;address segment 0000
01DD 1E PUSH DS ;save DS
01DE 8E D8 MOV DS,AX
01E0 A0 0417 MOV AL,DS:[417H] ;get shift/alternate data
01E3 1F POP DS
01E4 2E: 22 06 0115 R AND AL,CS:HMASK ;isolate alternate key
01E9 2E: 3A 06 0116 R CMP AL,CS:MKEY ;test for alternate key
01EE 74 2A JE VEC93 ;if hot-key found
01F0 VEC91:
01F0 51 PUSH CX ;add one to BCD COUNT
01F1 B9 0003 MOV CX,3
01F4 66| 2E: A1 010C R MOV EAX,CS:COUNT
01F9 66| 83 C0 01 ADD EAX,1
01FD 27 DAA ;make result BCD
01FE VEC92:
01FE 9C PUSHF
01FF 66| C1 C8 08 ROR EAX,8
0203 9D POPF
0204 14 00 ADC AL,0 ;propagate carry
0206 27 DAA
0207 E2 F5 LOOP VEC92
0209 66| C1 C8 08 ROR EAX,8
020D 66| 2E: A3 010C R MOV CS:COUNT,EAX
0212 59 POP CX
0213 66| 58 POP EAX
0215 2E: FF 2E 0108 R JMP CS:ADD9 ;do normal interrupt
021A VEC93:
021A FA CLI ;if hot-key pressed
021B E4 61 IN AL,61H ;interrupts off
021D 0C 80 OR AL,80H ;clear keyboard and
021F E6 61 OUT 61H,AL ;throw away hot key
0221 24 7F AND AL,7FH
0223 E6 61 OUT 61H,AL
0225 B0 20 MOV AL,20H ;reset keyboard interrupt
0227 E6 20 OUT 20H,AL
0229 FB STI ;enable interrupts
022A 2E: C6 06 0103 R MOV CS:HFLAG,1 ;indicate hot-key pressed
01 01
0230 66| 58 POP EAX
0232 CF IRET
0233 VEC9 ENDP
0233 VEC8 PROC FAR ;clock tick interrupt
0233 2E: 80 3E 0113 R CMP CS:FLAG8,0
00
0239 74 05 JZ VEC81 ;if not busy
023B 2E: FF 2E 0104 R JMP CS:ADD8 ;if busy
0240 VEC81:
0240 2E: 80 3E 0103 R CMP CS:HFLAG,0
00
0246 75 37 JNZ VEC83 ;if hot-key detected
0248 2E: 80 3E 0112 R CMP CS:SFLAG,0
00
024E 74 05 JZ VEC82 ;if start-up
0250 2E: FF 2E 0104 R JMP CS:ADD8 ;if not hot-key or start
0255 VEC82:

```

220 CHAPTER 7 PROGRAMMING THE MICROPROCESSOR

```

0255 9C                PUSHF                ;do old interrupt 8
0256 2E: FF 1E 0104 R  CALL CS:ADD8
025B 2E: C6 06 0113 R  MOV CS:FLAG8,1    ;indicate busy
01
0261 FB                STI                ;enable interrupts
0262 50                PUSH AX
0263 51                PUSH CX
0264 52                PUSH DX
0265 B4 02            MOV AH,2          ;get start-up time
0267 CD 1A            INT 1AH
0269 2E: 88 2E 0110 R  MOV CS:HOURL,CH   ;save hour
026E 2E: 88 0E 0111 R  MOV CS:MIN,CL     ;save minute
0273 5A                POP DX           ;restore registers
0274 59                POP CX
0275 58                POP AX
0276 2E: C6 06 0112 R  MOV CS:SFLAG,1   ;indicate started
01
027C E9 00A5          JMP VEC89         ;end it
027F                VEC83:                ;do hot-key display
027F 9C                PUSHF                ;do old interrupt 8
0280 2E: FF 1E 0104 R  CALL CS:ADD8
0285 2E: C6 06 0113 R  MOV CS:FLAG8,1    ;indicate busy
01
028B FB                STI                ;enable interrupts
028C 50                PUSH AX           ;save registers
028D 53                PUSH BX
028E B4 0F            MOV AH,0FH        ;get video mode
0290 CD 10            INT 10H
0292 3C 03            CMP AL,3
0294 76 05            JBE VEC84         ;if DOS text mode
0296 5B                POP BX           ;ignore if graphics mode
0297 58                POP AX
0298 E9 0083          JMP VEC88
029B                VEC84:                ;for text mode
029B 51                PUSH CX
029C 66 52            PUSH EDX
029E 57                PUSH DI
029F 56                PUSH SI
02A0 1E                PUSH DS
02A1 06                PUSH ES
02A2 FC                CLD
02A3 8C C8            MOV AX,CS         ;address this segment
02A5 8E C0            MOV ES,AX
02A7 B8 B800          MOV AX,0B800H    ;address text memory
02AA 8E D8            MOV DS,AX
02AC B9 00A0          MOV CX,160        ;save top screen line
02AF BF 0117 R        MOV DI,OFFSET SCRN
02B2 BE 0000          MOV SI,0
02B5 F3 A4            REP MOVSB
02B7 1E                PUSH DS           ;swap segments
02B8 06                PUSH ES
02B9 1F                POP DS
02BA 07                POP ES
02BB BF 0050          MOV DI,80         ;start display at center
02BE BE 01B7 R        MOV SI,OFFSET MES1
02C1 B4 0F            MOV AH,0FH        ;load white on black
02C3 B9 0007          MOV CX,7
02C6                VEC85:                ;display "Time = "
02C6 AC                LODSB
02C7 AB                STOSW
02C8 E2 FC            LOOP VEC85
02CA 2E: 8A 16 0111 R  MOV DL,CS:MIN
02CF 2E: 8A 36 0110 R  MOV DH,CS:HOURL

```

```

02D4 66| C1 E2 10      SHL  EDX,16
02D8 B9 0002          MOV  CX,2
02DB B3 30           MOV  BL,30H
02DD E8 004B        CALL DISP           ;display hours
02E0 B0 3A          MOV  AL,':'
02E2 AB             STOSW              ;display colon
02E3 B9 0002          MOV  CX,2
02E6 B3 80           MOV  BL,80H
02E8 E8 0040        CALL DISP           ;display minutes
02EB BE 01BE R      MOV  SI,OFFSET MES2 ;display KeyStrokes =
02EE B9 0010          MOV  CX,16
02F1                VEC86:
02F1 AC             LODSB
02F2 AB             STOSW
02F3 E2 FC          LOOP VEC86
02F5 66| 2E: 8B 16 010C R MOV  EDX,CS:COUNT ;get count
02FB B9 0008          MOV  CX,8
02FE B3 30           MOV  BL,30H
0300 E8 0028        CALL DISP           ;display count
0303                VEC87:
0303 B4 01           MOV  AH,1          ;wait for any key (BIOS)
0305 CD 16          INT  16H
0307 74 FA          JZ   VEC87
0309 FC             CLD
030A BE 0117 R      MOV  SI,OFFSET SCRN ;restore text
030D BF 0000          MOV  DI,0
0310 B9 00A0        MOV  CX,160
0313 F3/ A4         REP  MOVSB
0315 07             POP  ES
0316 1F             POP  DS
0317 5E             POP  SI
0318 5F             POP  DI
0319 66| 5A         POP  EDX
031B 59             POP  CX
031C 5B             POP  BX
031D 58             POP  AX
031E                VEC88:
031E 2E: C6 06 0103 R MOV  CS:HFLAG,0     ;kill hot-key
00
0324                VEC89:
0324 2E: C6 06 0113 R MOV  CS:FLAG8,0     ;indicate not busy
00
032A CF             IRET
032B                VEC8  ENDP
;
;The DISP procedure displays the BCD contents of EDX.
;***input parameters***
;CX = number of digits
;BL = 30H for blank leading zeros or 80H for no blanking
;ES = segment address of text mode display
;DI = offset address of text mode display
;
032B                DISP  PROC NEAR           ;display
032B 66| C1 C2 04      ROL  EDX,4          ;position number
032F 8A C2          MOV  AL,DL
0331 24 0F          AND  AL,0FH
0333 04 30          ADD  AL,30H         ;convert to ASCII
0335 AB             STOSW              ;store in text display
0336 3A C3          CMP  AL,BL         ;test for blanking
0338 74 04          JE   DISP1         ;if blanking needed
033A B3 80          MOV  BL,80H        ;turn off blanking

```

222 CHAPTER 7 PROGRAMMING THE MICROPROCESSOR

```

033C EB 03          JMP  DISP2          ;continue
033E                DISP1:
033E 83 EF 02      SUB  DI,2          ;blank digit
0341                DISP2:
0341 E2 E8        LOOP DISP
0343 C3           RET

0344                DISP  ENDP

0344                INSTALL:          ;install VEC8 and VEC9

0344 8C C8        MOV  AX,CS          ;load DS
0346 8E D8        MOV  DS,AX

0348 B8 3508      MOV  AX,3508H       ;get current vector 8
034B CD 21        INT  21H         ;and save it
034D 89 1E 0104 R MOV  WORD PTR ADD8,BX
0351 8C 06 0106 R MOV  WORD PTR ADD8+2,ES

0355 B8 3509      MOV  AX,3509H       ;get current vector 9
0358 CD 21        INT  21H         ;and save it
035A 89 1E 0108 R MOV  WORD PTR ADD9,BX
035E 8C 06 010A R MOV  WORD PTR ADD9+2,ES

0362 B8 2508      MOV  AX,2508H
0365 BA 0233 R    MOV  DX,OFFSET VEC8 ;address interrupt procedure
0368 CD 21        INT  21H         ;install vector 8

036A B8 2509      MOV  AX,2509H
036D BA 01CE R    MOV  DX,OFFSET VEC9 ;address interrupt procedure
0370 CD 21        INT  21H         ;install vector 9

0372 BA 0344 R    MOV  DX,OFFSET INSTALL ;find paragraphs
0375 C1 EA 04     SHR  DX,4
0378 42          INC  DX

0379 B8 3100      MOV  AX,3100H       ;set as a TSR
037C CD 21        INT  21H
                END

```

Note that the pop-up portion of this program only functions in the text mode and will count any unseen keystrokes that DOS generates. It also counts shift, alternate, and other keys as they are pressed and released. For example, the capital A will be counted as two or three keystrokes. This means that the count will be inflated. Even so, this program is useful for counting keystrokes by a given operator. If the operator reboots the system, the new reboot time is displayed and the count is cleared to zero.

The VEC9 interrupt service procedure intercepts all keystrokes. The IN AL,60H instruction reads the scan code from the keyboard interface within the personal computer. This is then tested for the K scan code. (Refer to Table 7-3 for the key scan codes.) If the K scan code is not found, the procedure increments the BCD count stored at location COUNT and returns to the normal keyboard interrupt handler. If the K scan code is detected, the contents of memory location 0000:0417 are tested for the alternate key. If an alternate key is detected, the program sets the HFLAG to 1, tosses away the hot key, and returns. Notice how the hot key is discarded by strobing I/O port number 61H. The keyboard is cleared by sending a logic 1 in bit position 7 of port 61H, followed by sending a logic 0 in bit position 7. The interrupt controller in the computer must also be cleared by sending a 20H out to I/O port number 20H.

The VEC8 interrupt service procedure tests the HFLAG for the hot key and the SFLAG for system startup. If the SFLAG = 0, the system has just been installed and the time is stored in HOUR and MIN. If the HFLAG = 1, a hot key was detected by VEC9. The VEC8 procedure responds to the hot key by storing the contents of the top line of the text display at memory array SCRN. Once the top line of the text display is stored, the message "Time =" is displayed, followed by the installation time. Next, the message "KeyStrokes =" is displayed, followed by the BCD number stored in COUNT. Recall that count is incremented each time VEC9 detects that a key is typed on the keyboard.

7-5 SUMMARY

1. The assembler program assembles modules that contain PUBLIC variables and segments, plus EXTRN (external) variables. The linker program links modules and library files to create a run-time program executed from the DOS command line. The run-time program usually has the extension EXE.
2. The MACRO and ENDM directives create a new opcode for use in programs. These macros are similar to procedures, except that there is no call or return. In place of them, the assembler inserts the code of the macro sequence into a program each time it is invoked. Macros can include variables that pass information and data to the macro sequence.
3. The DOS INT 21H function call provides a method of using the keyboard and video display. Function number 06H, placed into register AH, provides an interface to the keyboard and display. If DL = 0FFH, this function tests the keyboard for a keystroke. If no keystroke is detected, it returns equal. If a keystroke is detected, the standard ASCII character returns in AL. If an extended ASCII character is typed, it returns with AL = 00H, where the function must again be called to return with the extended ASCII character in AL. To display a character, DL is loaded with the character and AH with 06H before the INT 21H is used in a program.
4. Character strings are displayed by using function number 09H. The DS:DX register combination addresses the character string, which must end with a \$.
5. The INT 10H instruction accesses video BIOS (basic I/O system) procedures that control the video display and keyboard. The video BIOS functions are independent of DOS and function with any operating system.
6. The mouse driver is installed at interrupt vector 33H.
7. Data conversion from binary to BCD is accomplished with the AAM instruction for numbers that are less than 100 or by repeated division by 10 for larger numbers. Once converted to BCD, a 30H is added to convert each digit to ASCII code for the video display.
8. When converting from an ASCII number to BCD, a 30H is subtracted from each digit. To obtain the binary equivalent, we multiply by 10.
9. Lookup tables are used for code conversion with the XLAT instruction if the code is an eight-bit code. If the code is wider than eight bits, a short procedure that accesses a lookup table provides the conversion. Lookup tables are also used to hold addresses so that different parts of a program or different procedures can be selected.
10. Interrupt hooks allow application software to gain access to or intercept an interrupt. We often hook into the timer click interrupt (vector 8) or the keyboard interrupt (vector 9).
11. A terminate and stay resident (TSR) program is a program that remains in the memory and is often accessed through a hooked interrupt, using either the timer click or a hot key.
12. A hot key is a key that activates a terminate and stay resident program through the keyboard interrupt hook.

7-6 QUESTIONS AND PROBLEMS

1. The assembler converts a source file to a(n) _____ file.
2. What files are generated from the source file TEST.ASM if it is processed by MASM?
3. The linker program links object files and _____ files to create an execution file.
4. What does the PUBLIC directive indicate when placed in a program module?
5. What does the EXTRN directive indicate when placed in a program module?
6. What directives appear with labels defined external?

7. Describe how a library file works when it is linked to other object files by the linker program.
8. What assembler language directives delineate a macro sequence?
9. What is a macro sequence?
10. How are parameters transferred to a macro sequence?
11. Develop a macro called ADD32 that adds the 32-bit contents of DX–CX to the 32-bit contents of BX–AX.
12. How is the LOCAL directive used within a macro sequence?
13. Develop a macro called ADDLIST PARA1,PARA2 that adds the contents of PARA1 to PARA2. Each of these parameters represents an area of memory. The number of bytes added are indicated by register CX before the macro is invoked.
14. Develop a macro that sums a list of byte-sized data invoked by the macro ADDM LIST,LENGTH. The label LIST is the starting address of the data block and length is the number of data added. The result must be a 16-bit sum found in AX at the end of the macro sequence.
15. What is the purpose of the INCLUDE directive?
16. Develop a procedure called RANDOM. This procedure must return an eight-bit random number in register CL at the end of the subroutine. (One way to generate a random number is to increment CL each time the DOS function 06H tests the keyboard and finds *no* keystroke. In this way, a random number is generated.)
17. Develop a macro that uses the REPEAT statement to insert 10 NOP instructions in a program.
18. Develop a macro that uses the IFB/IFNB statements to test the parameter PARA in the macro DISP MACRO PARA. If PARA is blank, display a carriage return/line feed combination. If PARA is not blank, display PARA as an ASCII-coded character.
19. Develop a procedure that displays a character string that ends with a 00H. Your procedure must use the DS:DX register to address the start of the character string.
20. Develop a procedure that reads a key and displays the hexadecimal value of an extended ASCII-coded keyboard character if it is typed. If a normal character is typed, ignore it.
21. Use BIOS INT 10H to develop a procedure that positions the cursor at line 3, column 6.
22. When a number is converted from binary to BCD, the _____ instruction accomplishes the conversion, provided the number is less than 100 decimal.
23. How is a large number (over 100 decimal) converted from binary to BCD?
24. A BCD digit is converted to ASCII code by adding a(n) _____.
25. An ASCII-coded number is converted to BCD by subtracting _____.
26. Develop a procedure that reads an ASCII number from the keyboard and stores it as a BCD number into memory array DATA. The number ends when anything other than a number is typed.
27. Explain how a three-digit ASCII-coded number is converted to binary.
28. Develop a procedure that converts all lowercase ASCII-coded letters into uppercase ASCII-coded letters. Your procedure may not change any other character except the letters a–z.
29. Develop a lookup table that converts hexadecimal data 00H–0FH into the ASCII-coded characters that represent the hexadecimal digits. Make sure to show the lookup table and any software required for the conversion.
30. Develop a program sequence that jumps to memory location ONE if AL = 6, TWO if AL = 7, and THREE if AL = 8.
31. Show how to use the XLAT instruction to access a lookup table called LOOK that is located in the stack segment.
32. Develop a short sequence of instructions that place the line MOV AL,6 into a program if the contents memory location BED are true. You must use the IF statement.
33. Write a program that displays the binary powers of 2 (in decimal) on the video screen for the powers 0 through 7. Your display shows $2^n = \text{value}$ for each power of 2.

34. Using the technique discussed in question number 16, develop a program that displays random numbers between 1 and 47 (or whatever) for your state's lottery.
35. Develop a program the hooks into interrupt vector 0 to display the following message on a divide error: "Oops, you have attempted to divide by 0."

CHAPTER 8

8086/8088 Hardware Specifications

INTRODUCTION

In this chapter, we describe the pin functions of both the 8086 and 8088 microprocessors and provide details on the following hardware topics: clock generation, bus buffering, bus latching, timing, wait states, and minimum mode operation versus maximum mode operation. The simple microprocessors are explained first, because of their simple structures, as an introduction to the Intel microprocessor family.

Before it is possible to connect or interface anything to the microprocessor, it is necessary to understand the pin functions and timing. Thus, the information in this chapter is essential to a complete understanding of memory and I/O interfacing, which we cover in the later chapters of the text.

CHAPTER OBJECTIVES

Upon completion of this chapter, you will be able to:

1. Describe the function of each 8086 and 8088 pin.
2. Understand the microprocessor's DC characteristics and indicate its fan-out to common logic families.
3. Use the clock generator chip (8284A) to provide the clock for the microprocessor.
4. Connect buffers and latches to the buses.
5. Interpret the timing diagrams.
6. Describe wait states and connect the circuitry required to cause various numbers waits.
7. Explain the difference between minimum and maximum mode operation.

8-1 PIN-OUTS AND THE PIN FUNCTIONS

In this section, we explain the function and (in certain instances) the multiple functions of each of the microprocessor's pins. In addition, we discuss the DC characteristics to provide a basis for understanding the later sections on buffering and latching.

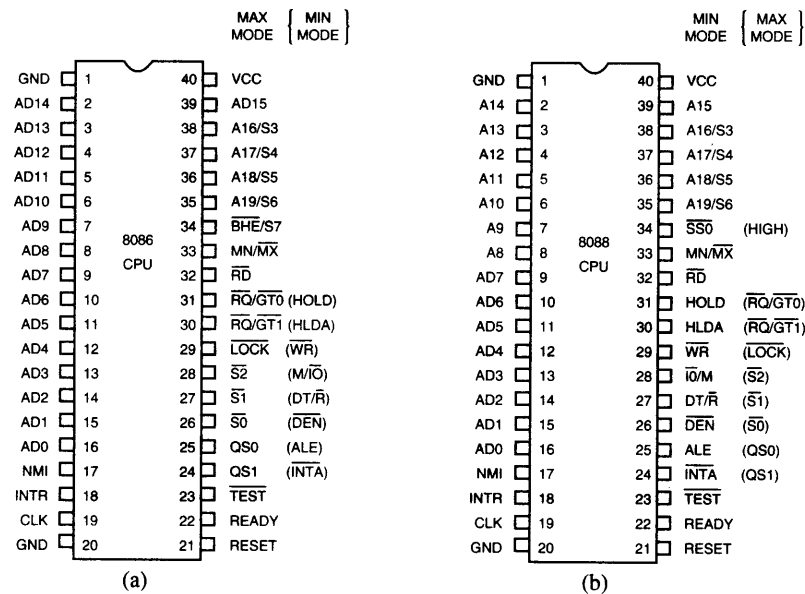


FIGURE 8-1 (a) The pin-out of the 8086 microprocessor; (b) the pin-out of the 8088 microprocessor.

The Pin-Out

Figure 8-1 illustrates the pin-outs of the 8086 and 8088 microprocessors. As a close comparison reveals, there is virtually no difference between these two microprocessors—both are packaged in 40-pin **dual in-line packages** (DIPs).

As mentioned in Chapter 1, the 8086 is a 16-bit microprocessor with a 16-bit data bus, and the 8088 is a 16-bit microprocessor with an 8-bit data bus. (As the pin-outs show, the 8086 has pin connections AD0–AD15, and the 8088 has pin connections AD0–AD7.) Data bus width is therefore the only major difference between these microprocessors.

There is, however, a minor difference in one of the control signals. The 8086 has an M/IO pin, and the 8088 has an IO/M pin. The only other hardware difference appears on Pin 34 of both chips: on the 8088, it is an SSO pin, while on the 8086, it is a BHE/S7 pin.

Power Supply Requirements

Both the 8086 and 8088 microprocessors require +5.0 V with a supply voltage tolerance of ± 10 percent. The 8086 uses a maximum supply current of 360 mA, and the 8088 draws a maximum of 340 mA. Both microprocessors operate in ambient temperatures of between 32° F and about 180° F. This range is not wide enough to be used outdoors in the winter or even in the summer, but extended temperature-range versions of the 8086 and 8088 microprocessors are available. There is also a CMOS version, which requires a very low supply current and has an extended temperature range. The 80C88 and 80C86 are CMOS versions that require only 10 mA of power supply current and function in temperature extremes of -40° F through $+225^{\circ}$ F.

DC Characteristics

It is impossible to connect anything to the pins of the microprocessor without knowing the input current requirement for an input pin and the output current drive capability for an output pin. This knowledge allows the hardware designer to select the proper interface components for use with the microprocessor without the fear of damaging anything.

TABLE 8-1 Input characteristics of the 8086 and 8088 microprocessors.

<i>Logic Level</i>	<i>Voltage</i>	<i>Current</i>
0	0.8 V maximum	±10 μA maximum
1	2.0 V minimum	±10 μA maximum

TABLE 8-2 Output characteristics of the 8086 and 8088 microprocessors.

<i>Logic Level</i>	<i>Voltage</i>	<i>Current</i>
0	0.45 V maximum	2.0 mA maximum
1	2.4 V minimum	-400 μA maximum

Input Characteristics. The input characteristics of these microprocessors are compatible with all the standard logic components available today. Table 8-1 depicts the input voltage levels and the input current requirements for any input pin on either microprocessor. The input current levels are very small because the inputs are the gates connections of MOSFETs and represent only leakage currents.

Output Characteristics. Table 8-2 illustrates the output characteristics of all the output pins of these microprocessors. The logic 1 voltage level of the 8086/8088 is compatible with that of most standard logic families, but the logic 0 level is not. Standard logic circuits have a maximum logic 0 voltage of 0.4 V, and the 8086/8088 has a maximum of 0.45 V. Thus, there is a difference of 0.05 V.

This difference reduces the noise immunity from a standard level of 400 mV (0.8 V - 0.45 V) to 350 mV. (The **noise immunity** is the difference between the logic 0 output voltage and the logic 0 input voltage levels.) This reduced noise immunity may result in problems with long wire connections or too many loads. It is therefore recommended that no more than 10 loads of any type or combination be connected to an output pin without buffering. If this loading factor is exceeded, noise will begin to take its toll in timing problems.

Table 8-3 lists some of the more common logic families and the recommended fan-out from the 8086/8088. The best choice of component types for the connection to an 8086/8088 output pin is a LS, 74ALS, or 74HC logic component. Note that some of the fan-out currents calculate to more than 10 unit loads. It is therefore recommended that if a fan-out of more than 10 unit loads is required, the system should be buffered.

Pin Connections

AD7-AD0

The 8088 **address/data bus** lines compose the multiplexed address data bus of the 8088 and contain the rightmost eight bits of the memory address or I/O port number whenever ALE is active (logic 1) or data whenever ALE is active (logic 0). These pins are at their high-impedance state during a hold acknowledge.

A15-A8

The 8088 **address bus** provides the upper-half memory address bits that are present throughout a bus cycle. These address connections go to their high-impedance state during a hold acknowledge.

AD15-AD8

The 8086 **address/data bus** lines compose the upper multiplexed address/data bus on the 8086. These lines contain address bits A15-A8 whenever ALE is a logic 1, and data bus connections D15-D8. These pins enter a high-impedance state whenever a hold acknowledge occurs.

TABLE 8-3 Recommended fan-out from any 8086/8088 pin connection.

Family	Sink Current	Source Current	Fan-out
TTL (74)	-1.6 mA	40 μ A	1
TTL (74LS)	-0.4 mA	20 μ A	5
TTL (74S)	-2.0 mA	50 μ A	1
TTL (74ALS)	-0.1 mA	20 μ A	10
TTL (74AS)	-0.5 mA	25 μ A	10
TTL (74F)	-0.5 mA	25 μ A	10
CMOS (74HC)	-10 μ A	10 μ A	10
CMOS (CD4)	-10 μ A	10 μ A	10
NMOS	-10 μ	10 μ A	10

A19/S6-A16/S3

The **address/status bus** bits are multiplexed to provide address signals A19-A16 and also status bits S6-S3. These pins also attain a high-impedance state during the hold acknowledge.

Status bit S6 always remains a logic 0, bit S5 indicates the condition of the IF flag bits, and S4 and S3 show which segment is accessed during the current bus cycle. See Table 8-4 for the truth table of S4 and S3. These two status bits can be used to address four separate 1M byte memory banks by decoding them as A21 and A20.

RD

Whenever the **read signal** is a logic 0, the data bus is receptive to data from the memory or I/O devices connected to the system. This pin floats to its high-impedance state during a hold acknowledge.

READY

This input is controlled to insert wait states into the timing of the microprocessor. If the READY pin is placed at a logic 0 level, the microprocessor enters into wait states and remains idle. If the READY pin is placed at a logic 1 level, it has no effect on the operation of the microprocessor.

INTR

Interrupt request is used to request a hardware interrupt. If INTR is held high when IF = 1, the 8086/8088 enters an interrupt acknowledge cycle (INTA becomes active) after the current instruction has complete execution.

TEST

The **Test** pin is an input that is tested by the WAIT instruction. If TEST is a logic 0, the WAIT instruction functions as a NOP. If TEST is a logic 1, the WAIT instruction waits for TEST to become a logic 0. This pin is most often connected to the 8087 numeric coprocessor.

NMI

The **non-maskable interrupt** input is similar to INTR except that the NMI interrupt does not check to see whether the IF flag bit is a logic 1. If NMI is activated, this interrupt input uses interrupt vector 2.

RESET

The **reset** input causes the microprocessor to reset itself if this pin is held high for a minimum of four clocking periods. Whenever the 8086 or 8088 is reset, it begins executing instructions at memory location FFFF0H and disables future interrupts by clearing the IF flag bit.

TABLE 8-4 Function of status bits S3 and S4.

S4	S3	Function
0	0	Extra segment
0	1	Stack segment
1	0	Code or no segment
1	1	Data segment

CLK	The clock pin provides the basic timing signal to the microprocessor. The clock signal must have a duty cycle of 33 percent (high for one-third of the clocking period and low for two-thirds) to provide proper internal timing for the 8086/8088.
VCC	This power supply input provides a +5.0 V, $\pm 10\%$ signal to the microprocessor.
GND	The ground connection is the return for the power supply. Note that the 8086/8088 microprocessors have two pins labeled GND—both must be connected to ground for proper operation.
MN/MX	The minimum/maximum mode pin selects either minimum mode or maximum mode operation for the microprocessor. If minimum mode is selected, the MN/MX pin must be connected directly to +5.0 V.
BHE/S7	The bus high enable pin is used in the 8086 to enable the most-significant data bus bits (D15–D8) during a read or a write operation. The state of S7 is always a logic 1.

Minimum Mode Pins. Minimum mode operation of the 8086/8088 is obtained by connecting the MN/MX pin directly to +5.0 V. Do not connect this pin to +5.0 V through a pull-up resistor or it will not function correctly.

IO/M or M/IO	The IO/M (8088) or the M/IO (8086) pin selects memory or I/O. This pin indicates that the microprocessor address bus contains either a memory address or an I/O port address. This pin is at its high-impedance state during a hold acknowledge.
WR	The write line is a strobe that indicates that the 8086/8088 is outputting data to a memory or I/O device. During the time that the WR is a logic 0, the data bus contains valid data for memory or I/O. This pin floats to a high-impedance during a hold acknowledge.
INTA	The interrupt acknowledge signal is a response to the INTR input pin. The INTA pin is normally used to gate the interrupt vector number onto the data bus in response to an interrupt request.
ALE	Address latch enable shows that the 8086/8088 address/data bus contains address information. This address can be a memory address or an I/O port number. Note that the ALE signal does not float during a hold acknowledge.
DT/R	The data transmit/receive signal shows that the microprocessor data bus is transmitting (DT/R = 1) or receiving (DT/R = 0) data. This signal is used to enable external data bus buffers.
DEN	Data bus enable activates external data bus buffers.
HOLD	The hold input requests a direct memory access (DMA). If the HOLD signal is a logic 1, the microprocessor stops executing software and places its address, data, and control bus at the high-impedance state. If the HOLD pin is a logic 0, the microprocessor executes software normally.
HLDA SS0	Hold acknowledge indicates that the 8086/8088 has entered the hold state The SS0 status line is equivalent to the S0 pin in maximum mode operation of the microprocessor. This signal is combined with IO/M and DT/R to decode the function of the current bus cycle (see Table 8–5)

Maximum Mode Pins. In order to achieve maximum mode for use with external coprocessors, connect the MN/MX pin to ground.

S2, S1, and S0 The status bits indicate the function of the current bus cycle. These signals are normally decoded by the 8288 bus controller described later in this chapter. Table 8–6 shows the function of these three status bits in the maximum mode.

TABLE 8-5 Bus cycle status (8088) using SS0.

<i>IO/M</i>	<i>DT/R</i>	<i>SS0</i>	<i>Function</i>
0	0	0	Interrupt acknowledge
0	0	1	Memory read
0	1	0	Memory write
0	1	1	Halt
1	0	0	Opcode fetch
1	0	1	I/O read
1	1	0	I/O write
1	1	1	Passive

TABLE 8-6 Bus control functions generated by the bus controller (8288) using S2, S1, and S0.

<i>S2</i>	<i>S1</i>	<i>S0</i>	<i>Function</i>
0	0	0	Interrupt acknowledge
0	0	1	I/O read
0	1	0	I/O write
0	1	1	Halt
1	0	0	Opcode fetch
1	0	1	Memory read
1	1	0	Memory write
1	1	1	Passive

TABLE 8-7 Queue status bits.

<i>QS1</i>	<i>QS0</i>	<i>Function</i>
0	0	Queue is idle
0	1	First byte of opcode
1	0	Queue is empty
1	1	Subsequent byte of opcode

RO/GT1 and RO/GT0

The **request/grant** pins request direct memory accesses (DMA) during maximum mode operation. These lines are bi-directional, and are used to both request and grant a DMA operation.

LOCK

The **lock output** is used to lock peripherals off the system. This pin is activated by using the LOCK: prefix on any instruction.

QS1 and QS0

The **queue status** bits show the status of the internal instruction queue. These pins are provided for access by the numeric coprocessor (8087). See Table 8-7 for the operation of the queue status bits.

8-2 CLOCK GENERATOR (8284A)

This section describes the clock generator (8284A), the RESET signal, and introduces the READY signal for the 8086/8088 microprocessors. (The READY signal and its associated circuitry are treated in detail in Section 8-5.)

The 8284A Clock Generator

The 8284A is an ancillary component to the 8086/8088 microprocessors. Without the clock generator, many additional circuits are required to generate the clock (CLK) in an 8086/8088-based system. The 8284A provides the

following basic functions or signals: clock generation, RESET synchronization, READY synchronization, and a TTL-level peripheral clock signal. Figure 8-2 illustrates the pin-out of the 8284A clock generator.

Pin Functions. The 8284A is an 18-pin integrated circuit, designed specifically for use with the 8086/8088 microprocessors. The following is a list of each pin and its function:

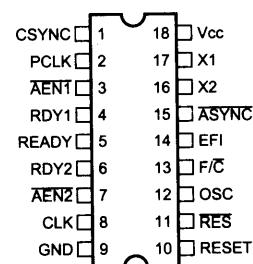


FIGURE 8-2 The pin-out of the 8284A clock generator.

AEN1 and AEN2	The address enable pins are provided to qualify the bus ready signals, RDY1 and RDY2, respectively. Section 8-5 illustrates the use of these two pins, which are used to cause wait states, along with the RDY1 and RDY2 inputs. Wait states are generated by the READY pin of the 8086/8088 microprocessors, which is controlled by these two inputs.
RDY1 and RDY2	The bus ready inputs are provided, in conjunction with the AEN1 and AEN2 pins, to cause wait states in an 8086/8088-based system.
ASYNC	The ready synchronization selection input selects either one or two stages of synchronization for the RDY1 and RDY2 inputs.
READY	Ready is an output pin that connects to the 8086/8088 READY input. This signal is synchronized with the RDY1 and RDY2 inputs.
X1 and X2	The crystal oscillator pins connect to an external crystal used as the timing source for the clock generator and all its functions.
F/C	The frequency/crystal select input chooses the clocking source for the 8284A. If this pin is held high, an external clock is provided to the EFI input pin; if it is held low, the internal crystal oscillator provides the timing signal.
EFI	The external frequency input is used when the F/C pin is pulled high. EFI supplies the timing whenever the F/C pin is high.
CLK	The clock output pin provides the CLK input signal to the 8086/8088 microprocessors and other components in the system. The CLK pin has an output signal that is one-third of the crystal or EFI input frequency, and has a 33-percent duty cycle, which is required by the 8086/8088.
PCLK	The peripheral clock signal is one-sixth the crystal or EFI input frequency, and has a 50-percent duty cycle. The PCLK output provides a clock signal to the peripheral equipment in the system.
OSC	The oscillator output is a TTL-level signal that is at the same frequency as the crystal or EFI input. The OSC output provides an EFI input to other 8284A clock generators in some multiple-processor systems.
RES	The reset input is an active-low input to the 8284A. The RES pin is often connected to an RC network that provides power-on resetting.
RESET	The reset output is connected to the 8086/8088 RESET input pin.
CSYNC	The clock synchronization pin is used whenever the EFI input provides synchronization in systems with multiple processors. If the internal crystal oscillator is used, this pin must be grounded.
GND	The ground pin connects to ground.
Vcc	This power supply pin connects to +5.0 V with a tolerance of ± 10 percent.

Operation of the 8284A

The 8284A is a relatively easy component to understand. Figure 8-3 illustrates the internal block diagram of the 8284A clock generator.

Operation of the Clock Section. The top half of the logic diagram represents the clock and reset synchronization section of the 8284A clock generator. As the diagram shows, the crystal oscillator has two inputs: X1 and X2. If a crystal is attached to X1 and X2, the oscillator generates a square-wave signal at the same frequency as the crystal. The square-wave signal is fed to an AND gate and also to an inverting buffer that provides the OSC output signal. The OSC signal is sometimes used as an EFI input to other 8284A circuits in a system.

An inspection of the AND gate reveals that when F/C is a logic 0, the oscillator output is steered through to the divide-by-3 counter. If F/C is a logic 1, then EFI is steered through to the counter.

The output of the divide-by-3 counter generates the timing for ready synchronization, a signal for another counter (divide-by-2), and the CLK signal to the 8086/8088 microprocessors. The CLK signal is also buffered before it leaves the clock generator. Notice that the output of the first counter feeds the second. These two cascaded counters provide the divide-by-6 output at PCLK, the peripheral clock output.

Figure 8-4 shows how an 8284A is connected to the 8086/8088. Notice that F/C and CSYNC are grounded to select the crystal oscillator; and that a 15 MHz crystal provides the normal 5 MHz clock signal to the 8086/8088, as well as a 2.5 MHz peripheral clock signal.

Operation of the Reset Section. The reset section of the 8284A is very simple: It consists of a Schmitt trigger buffer and a single D-type flip-flop circuit. The D-type flip-flop ensures that the timing requirements of the 8086/8088 RESET input are met. This circuit applies the RESET signal to the microprocessor on the negative edge of each clock. The 8086/8088 microprocessors sample RESET at the positive edge (1-to-0 transition) of each clock. The 8086/8088 microprocessors sample RESET at the positive edge (0-to-1 transition) of the clocks; therefore, this circuit meets the timing requirements of the 8086/8088.

Refer to Figure 8-4. Notice that an RC circuit provides a logic 0 to the RES input pin when power is first applied to the system. After a short time, the RES input becomes a logic 1 because the capacitor charges toward +5.0 V through the resistor. A push-button switch allows the microprocessor to be reset by the operator. Correct reset

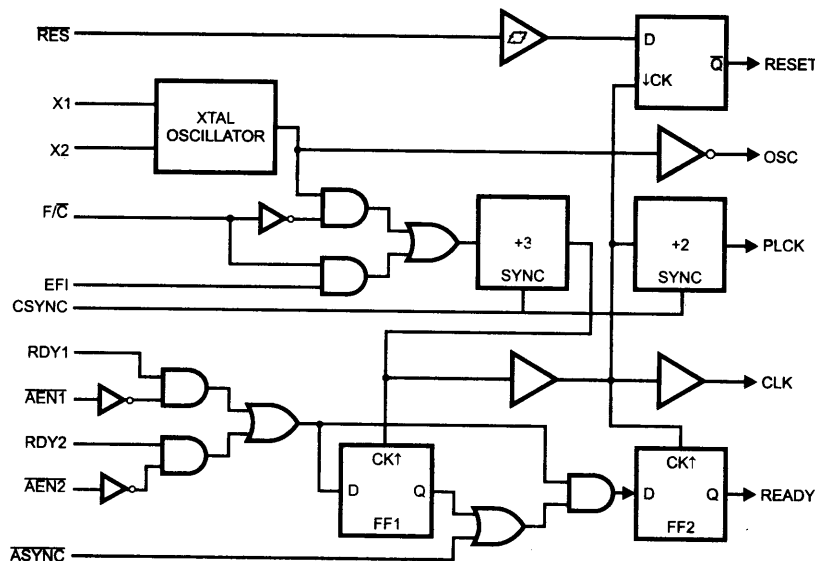


FIGURE 8-3 The internal block diagram of the 8284A clock generator.

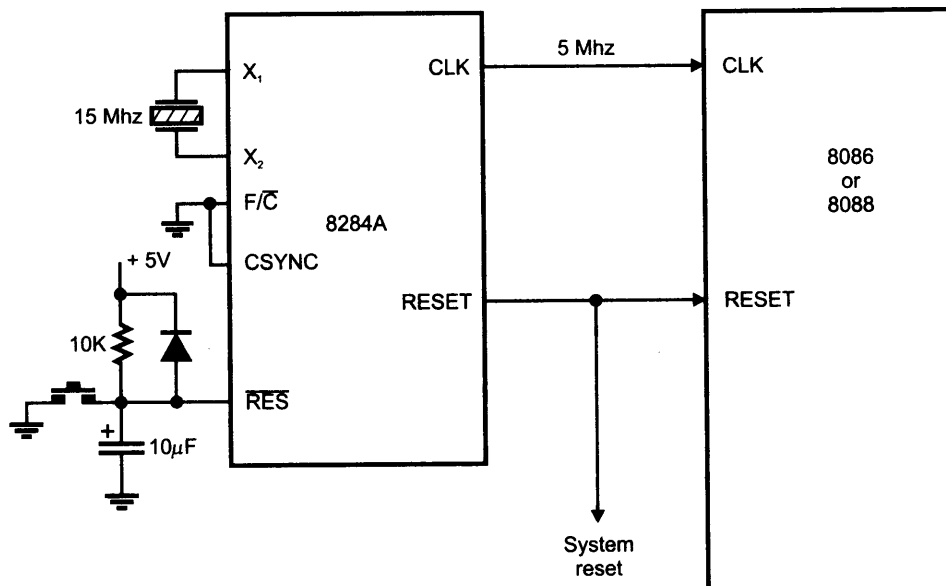


FIGURE 8-4 The clock generator (8284A) and the 8086 and 8088 microprocessor illustrating the connection for the clock and reset signals. A 15 MHz crystal provides the 5 MHz clock for the microprocessor.

timing requires the RESET input to become a logic 1 no later than four clocks after system power is applied, and to be held high for at least 50 μ s. The flip-flop makes certain that RESET goes high in four clocks, and the RC time constant ensures that it stays high for at least 50 μ s.

8-3 BUS BUFFERING AND LATCHING

Before the 8086/8088 microprocessors can be used with memory or I/O interfaces, their multiplexed buses must be demultiplexed. This section provides the detail required to demultiplex the buses and illustrates how the buses are buffered for very large systems. (Because the maximum fan-out is 10, the system must be buffered if it contains more than 10 other components.)

Demultiplexing the Buses

The address/data bus on the 8086/8088 is multiplexed (**shared**) to reduce the number of pins required for the 8086/8088 microprocessor integrated circuit. Unfortunately, this burdens the hardware designer with the task of extracting or demultiplexing information from these multiplexed pins.

Why not leave the buses multiplexed? Memory and I/O require that the address remains valid and stable throughout a read or write cycle. If the buses are multiplexed, the address changes at the memory and I/O, which causes them to read or write data in the wrong locations.

All computer systems have three buses: (1) an address bus that provides the memory and I/O with the memory address or the I/O port number, (2) a data bus that transfers data between the microprocessor and the memory and I/O in the system, and (3) a control bus that provides control signals to the memory and I/O. These buses must be present in order to interface to memory and I/O.

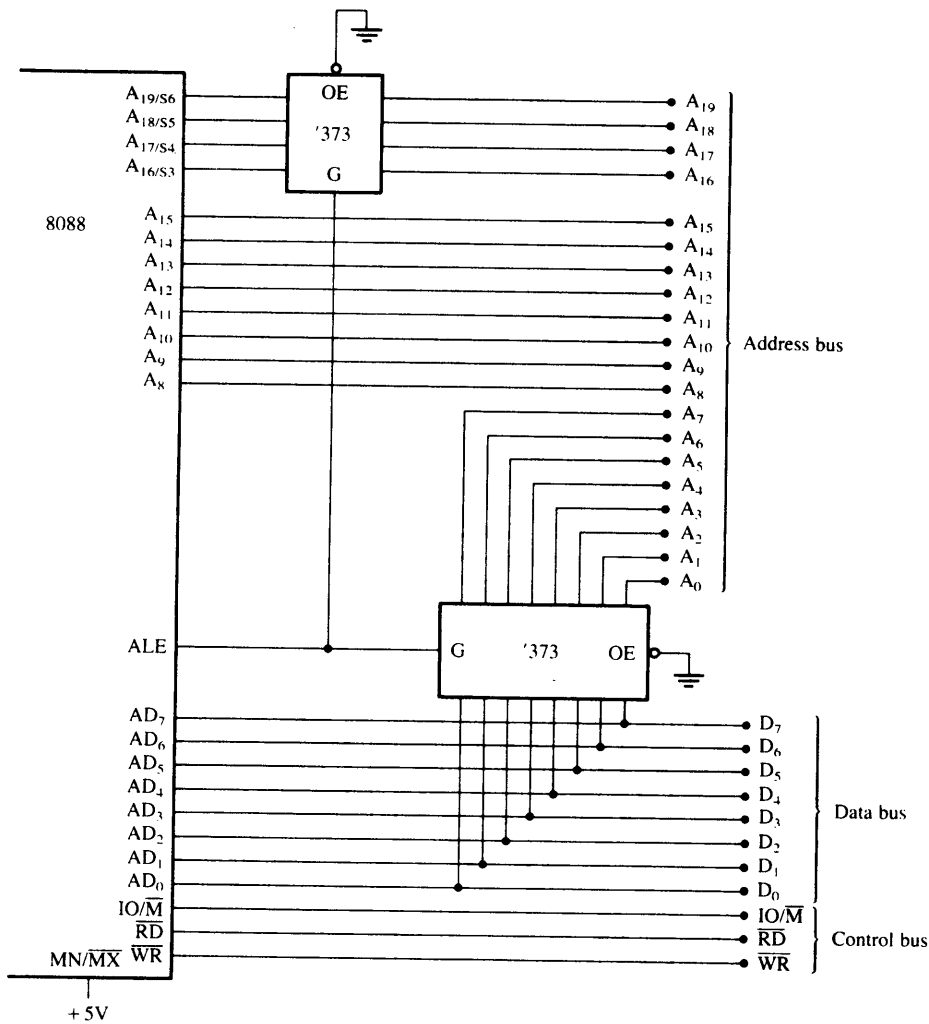


FIGURE 8-5 The 8088 microprocessor shown with a demultiplexed address bus. This is the model used to build many 8088-based systems.

Demultiplexing the 8088. Figure 8-5 illustrates the 8088 microprocessor and the components required to demultiplex its buses. In this case, two 74LS373 transparent latches are used to demultiplex the address/data bus connections AD7-AD0 and the multiplexed address/status connections A19/S6-A16/S3.

These transparent latches, which are like wires whenever the address latch enable pin (ALE) becomes a logic 1, pass the inputs to the outputs. After a short time, ALE returns to its logic 0 condition, which causes the latches to remember the inputs at the time of the change to a logic 0. In this case, A7-A0 are stored in the bottom latch and A19-A16 are stored in the top latch. This yields a separate address bus with connections A19-A0. These address connections allow the 8088 to address 1M bytes of memory space. The fact that the data bus is separate allows it to be connected to any eight-bit peripheral device or memory component.

Demultiplexing the 8086. Like the 8088, the 8086 system requires separate address, data, and control buses. It differs primarily in the number of multiplexed pins. In the 8088, only AD7-AD0 and A19/S6-A16/S3 are

multiplexed. In the 8086, the multiplexed pins include AD15–AD0, A19/S6–A16/S3, and $\overline{\text{BHE}}/\text{S7}$. All of these signals must be demultiplexed.

Figure 8–6 illustrates a demultiplexed 8086 with all three buses: address (A19–A0 and $\overline{\text{BHE}}$), data (D15–D0), and control (M/I $\overline{\text{O}}$, $\overline{\text{RD}}$, and $\overline{\text{WR}}$).

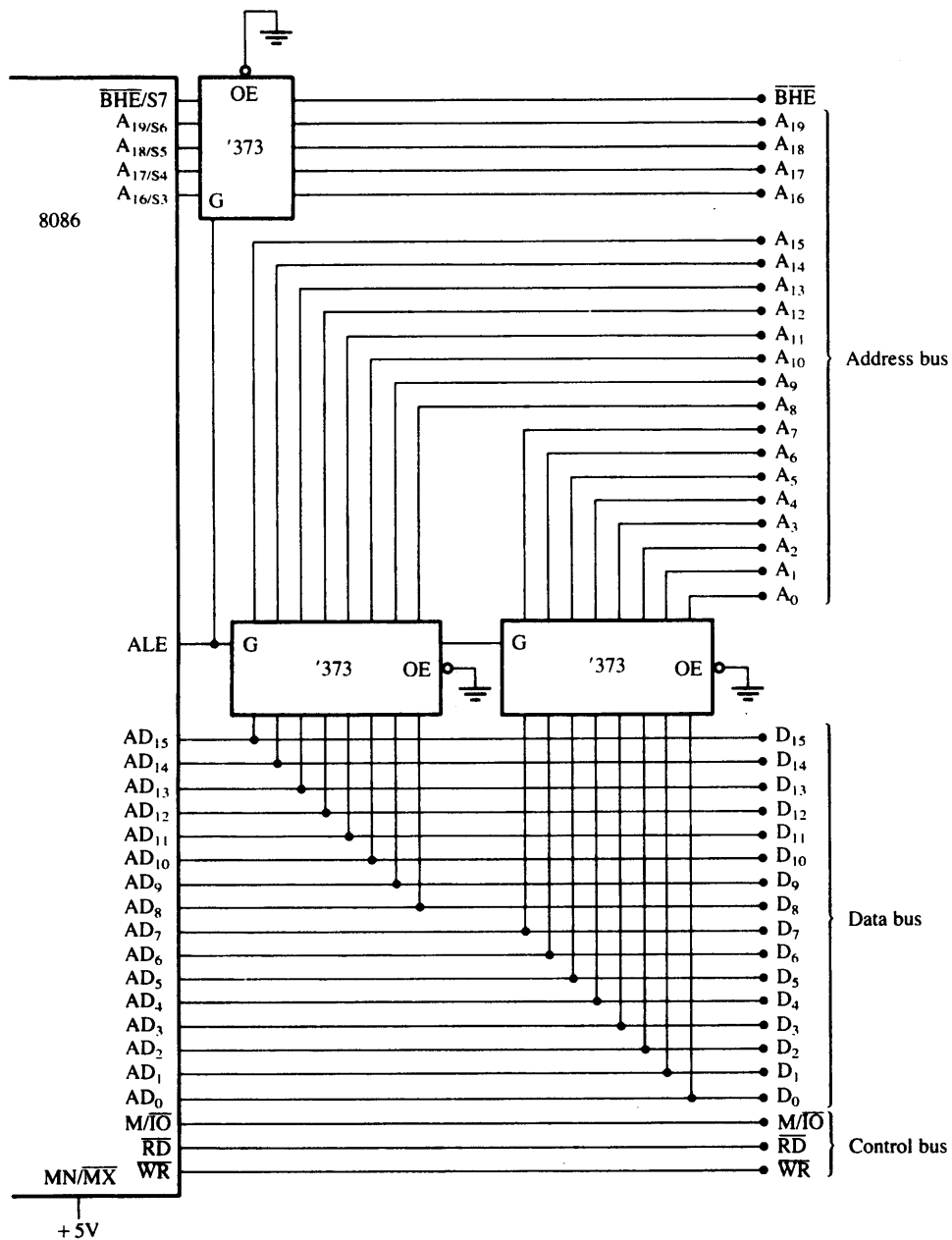


FIGURE 8–6 The 8086 microprocessor shown with a demultiplexed address bus. This is the model used to build many 8086-based systems.

This circuit shown in Figure 8-6 is almost identical to the one pictured in Figure 8-5, except that an additional 74LS373 latch has been added to demultiplex the address/data bus pins AD15-AD8 and a BHE/S7 input has been added to the top 74LS373 to select the high-order memory bank in the 16-bit memory system of the 8086. Here, the memory and I/O system see the 8086 as a device with a 20-bit address bus (A19-A0), a 16-bit data bus (D15-D0), and a three-line control bus (M/I/O), RD, and WR).

The Buffered System

If more than 10 unit loads are attached to any bus pin, the entire 8086 or 8088 system must be buffered. The demultiplexed pins are already buffered by the 74LS373 latches, which have been designed to drive the high-capacitance buses encountered in microcomputer systems. The buffer's output currents have been increased so that more TTL unit loads may be driven: a logic 0 output provides up to 32 mA of sink current, and a logic 1 output provides up to 5.2 mA of source current.

A fully buffered signal will introduce a timing delay to the system. This causes no difficulty unless memory or I/O devices are used, which function at near the maximum speed of the bus. Section 8-4 discusses this problem and the time delays involved in more detail.

The Fully Buffered 8088. Figure 8-7 depicts a fully buffered 8088 microprocessor. Notice that the remaining eight address pins, A15-A8, use a 74LS244 octal buffer; the eight data bus pins, D7-D0, use a 74LS245 octal bi-directional bus buffer; and the control bus signals, IO/M, RD, and WR, use a 74LS244 buffer. A fully-buffered 8088 system requires two 74LS244s, one 74LS245, and two 74LS373s. The direction of the 74LS245 is controlled by the DT/R signal, and is enabled and disabled by the DEN signal.

The Fully Buffered 8086. Figure 8-8 illustrates a fully buffered 8086 microprocessor. Its address pins are already buffered by the 74LS373 address latches; its data bus employs two 74LS245 octal bi-directional bus buffers; and the control bus signals, M/I/O, RD, and WR, use a 74LS244 buffer. A fully buffered 8086 system requires one 74LS244, two 74LS245s, and three 74LS373s. The 8086 requires one more buffer than the 8088 because of the extra eight data bus connections, D15-D8. It also has a BHE signal that is buffered for memory-bank selection.

8-4 BUS TIMING

It is essential to understand system bus timing before choosing a memory or I/O device for interfacing to the 8086 or 8088 microprocessors. This section provides insight into the operation of the bus signals, and the basic read and write timing of the 8086/8088. It is important to note that we discuss only the times that affect memory and I/O interfacing in this section.

Basic Bus Operation

The three buses of the 8086 and 8088—address, data, and control—function exactly the same way as those of any other microprocessor. If data are written to the memory (see the simplified timing for write in Figure 8-9), the microprocessor outputs the memory address on the address bus, outputs the data to be written into memory on the data bus, and issues a write (WR) to memory and IO/M = 0 for the 8088 and M/I/O = 1 for the 8086. If data are read from the memory (see the simplified timing for read in Figure 8-10), the microprocessor outputs the memory address on the address bus, issues a read (RD) memory signal, and accepts the data via the data bus.

Timing in General

The 8086/8088 microprocessors use the memory and I/O in periods called **bus cycles**. Each bus cycle equals four system-clocking periods (T states). Some new microprocessors divide the bus cycle into as few as two clocking

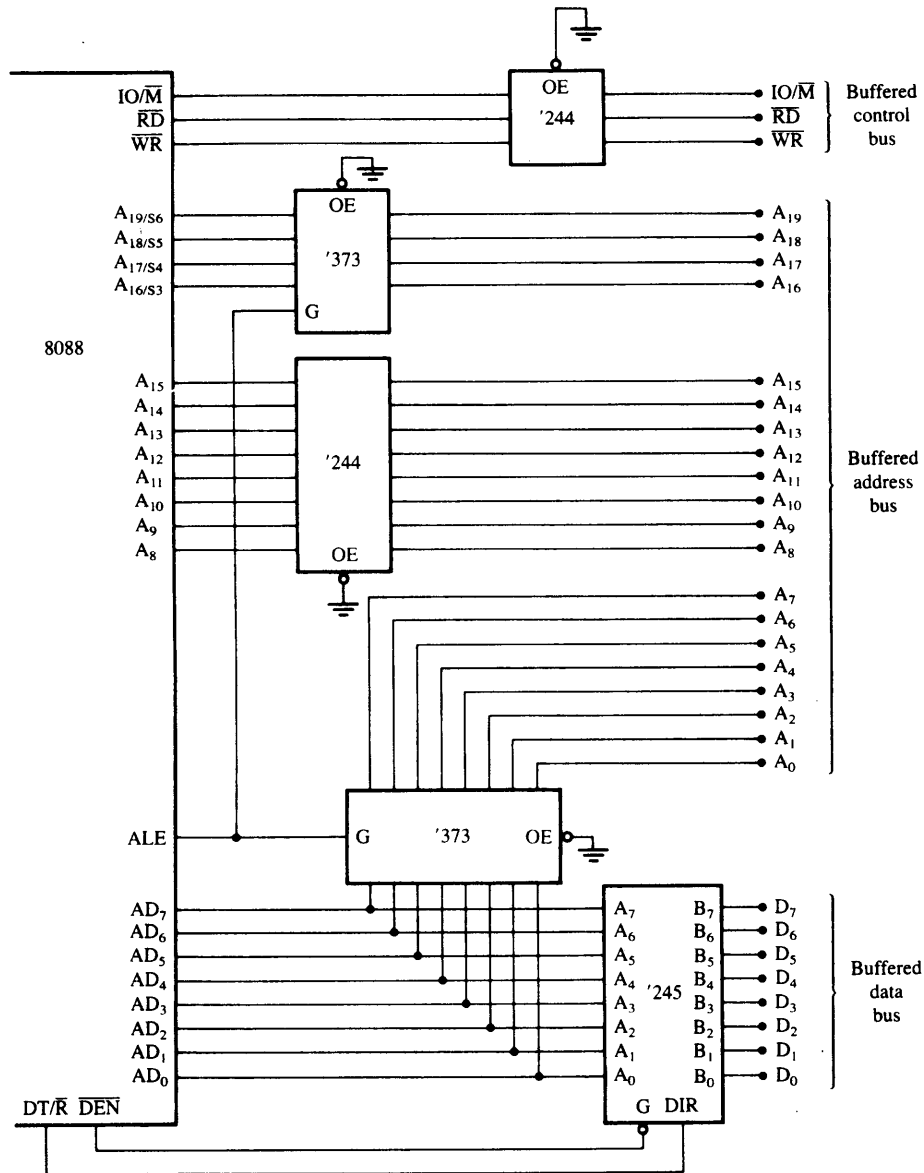


FIGURE 8-7 A fully buffered 8088 microprocessor.

periods. If the clock is operated at 5 MHz (the basic operating frequency for these two microprocessors), one 8086/8088 bus cycle is complete in 800 ns. This means that the microprocessor reads or writes data between itself and memory or I/O at a maximum rate of 1.25 million times a second. (Because of the internal queue, the 8086/8088 can execute 2.5 million instructions per second [MIPS] in bursts.) Other available versions of these microprocessors operate at much higher transfer rates due to higher clock frequencies.

During the first clocking period in a bus cycle, which is called T1, many things happen. The address of the memory or I/O location is sent out via the address bus and the address/data bus connections. (The address/data bus is multiplexed and sometimes contains memory-addressing information, sometimes data.) During T1, control